

- Differential Cryptanalysis of Simplified DES

- UCSB CS290g



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Outline

- What is Cryptanalysis
- Differential Cryptanalysis(DC)
- Simplified DES (SDES)
- DC of SDES
- Capturing the Key
- Conclusion

Cryptanalysis

What is Cryptanalysis

It is the analysis of a cryptographic system to find the plain-text. Some information is required to perform the cryptanalysis.

Differential Cryptanalysis

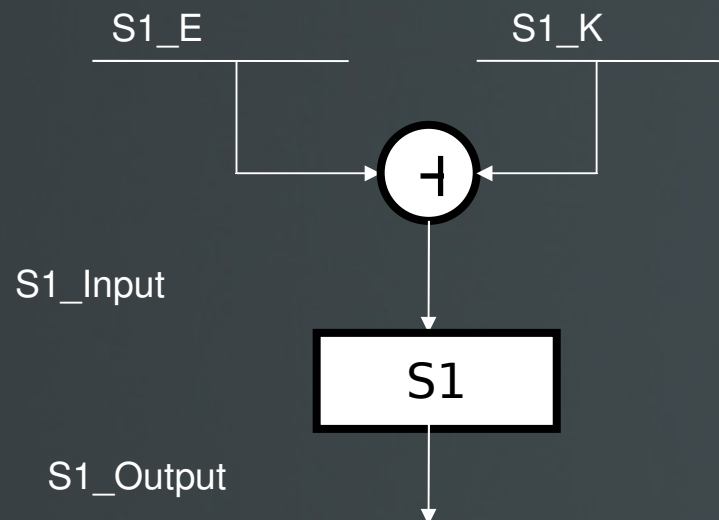
Study of the impact of difference in inputs on the difference in outputs

- DES does not exhibit pseudo randomness in differential output
- Key does not have any influence on the differential output

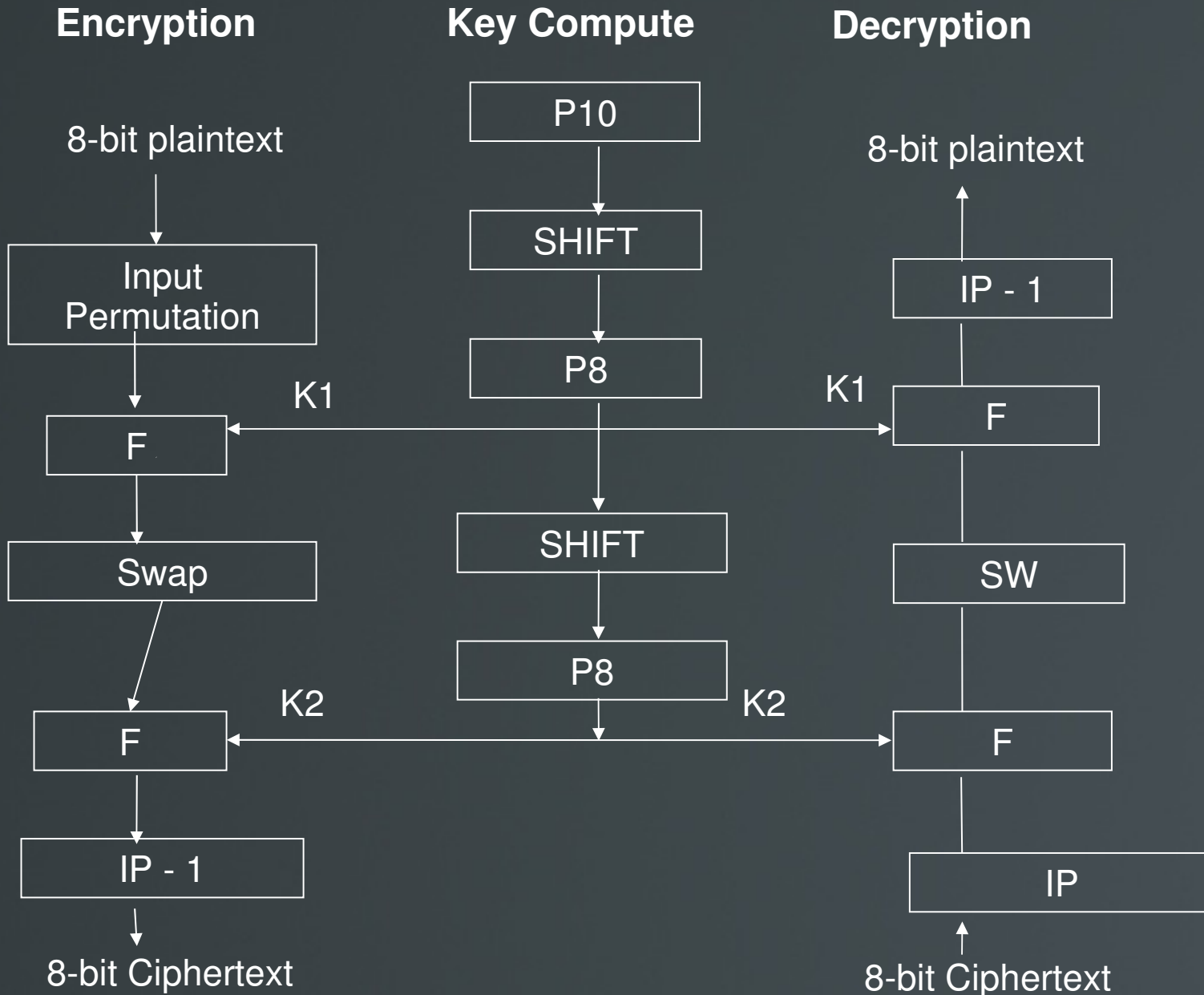
Differential Cryptanalysis

The Differential Input to the S-box remains same regardless of what the Key is

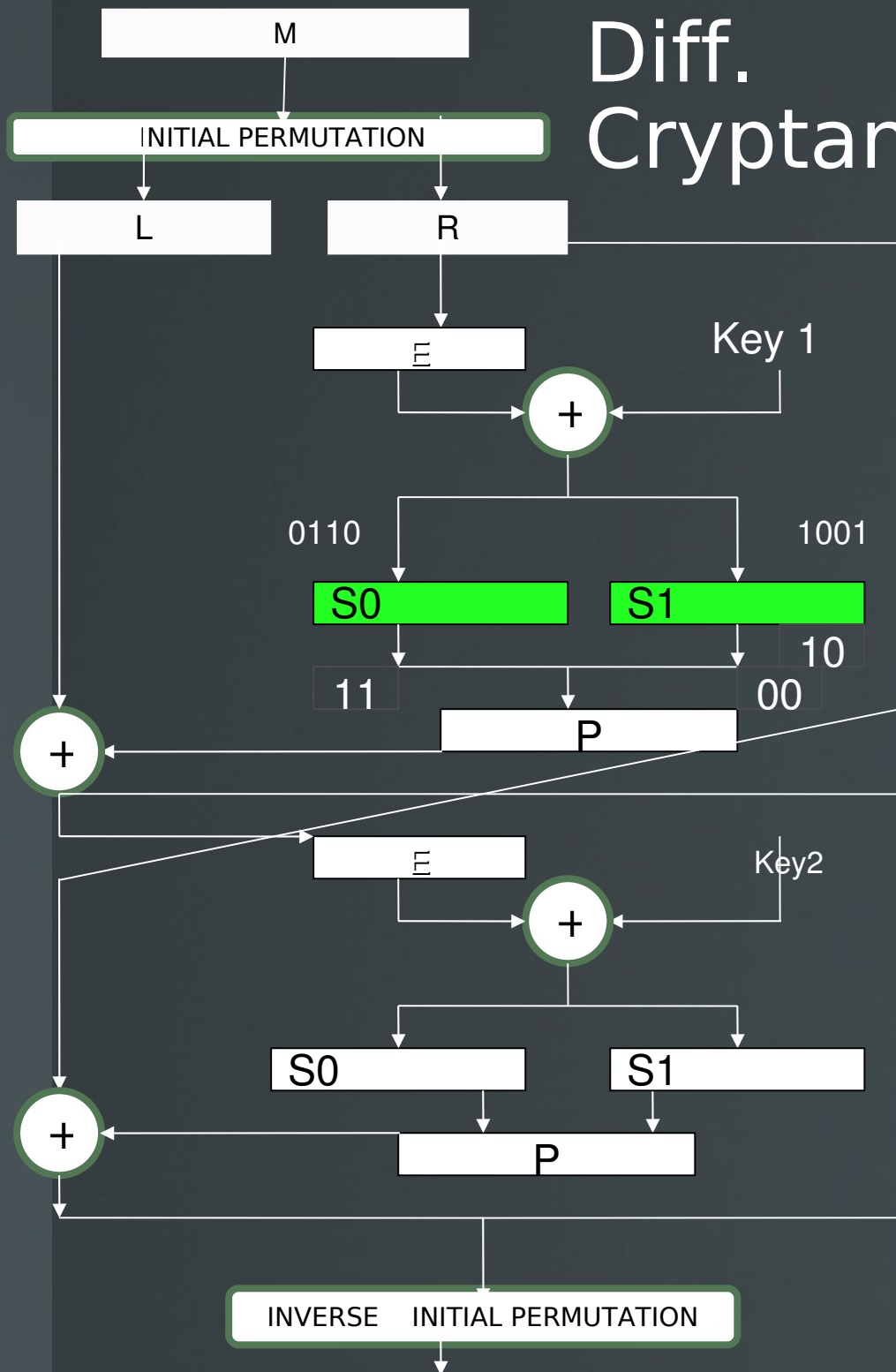
$$\begin{aligned} S1_Input &= S1_Input \oplus S1_K \\ &= (S1_E \oplus S1_K) \oplus (S1_K \oplus S1_K) \\ &= S1_E \oplus S1_K \oplus S1_K \\ &= S1_E \end{aligned}$$



Simplified DES



Diff. Cryptanal



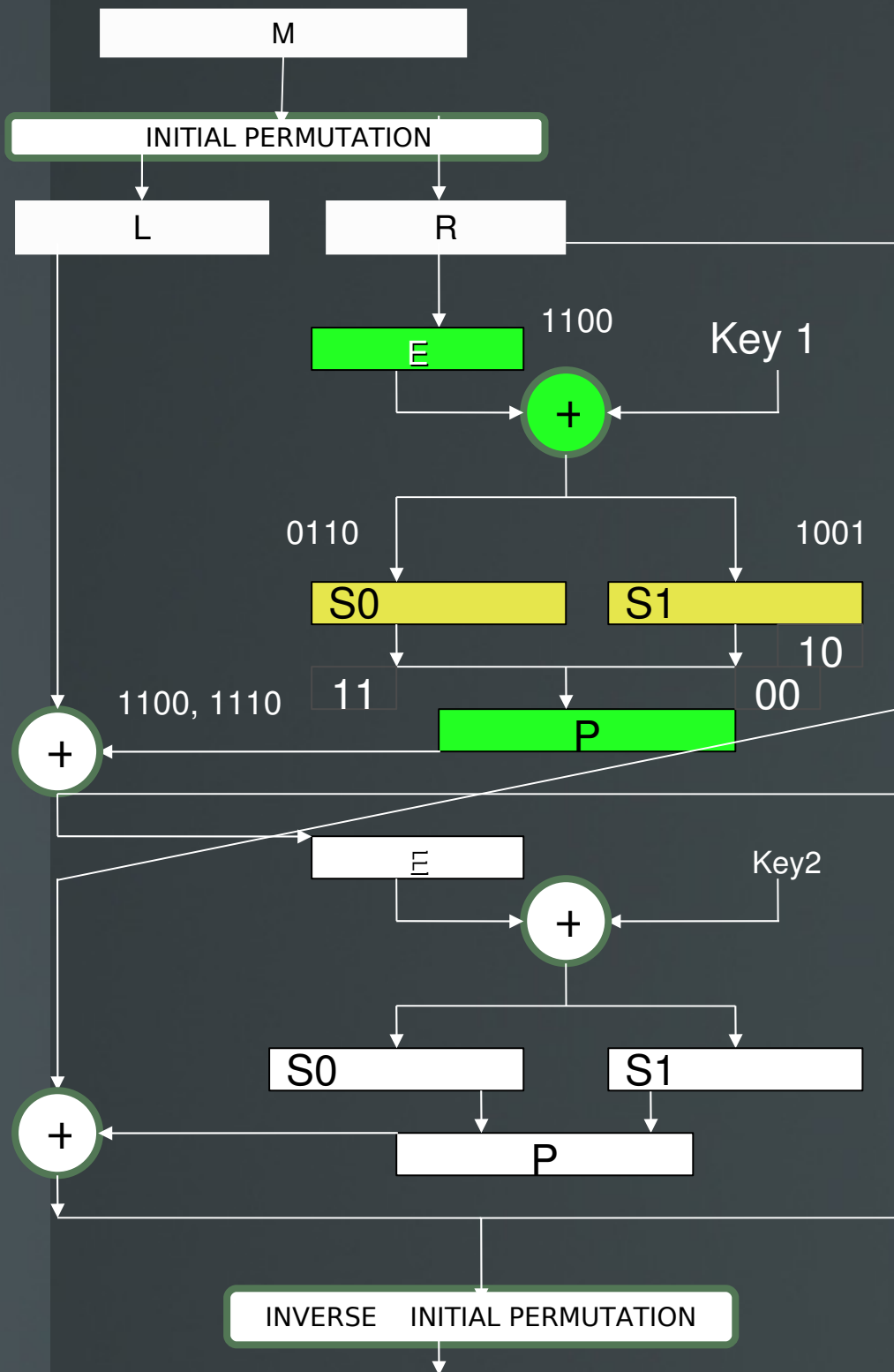
S0

<i>DX/DY</i>	0	1	2	3
0	16	0	0	0
1	0	2	10	4
2	0	10	6	0
3	2	4	0	10
4	2	6	6	2
5	10	0	4	2
6	0	2	2	12
7	4	10	2	0
8	2	4	8	2
9	8	2	2	4
10	4	2	2	8
11	2	8	4	2
12	8	2	2	4
13	1	4	8	2
14	1	8	4	2
15	4	2	2	8

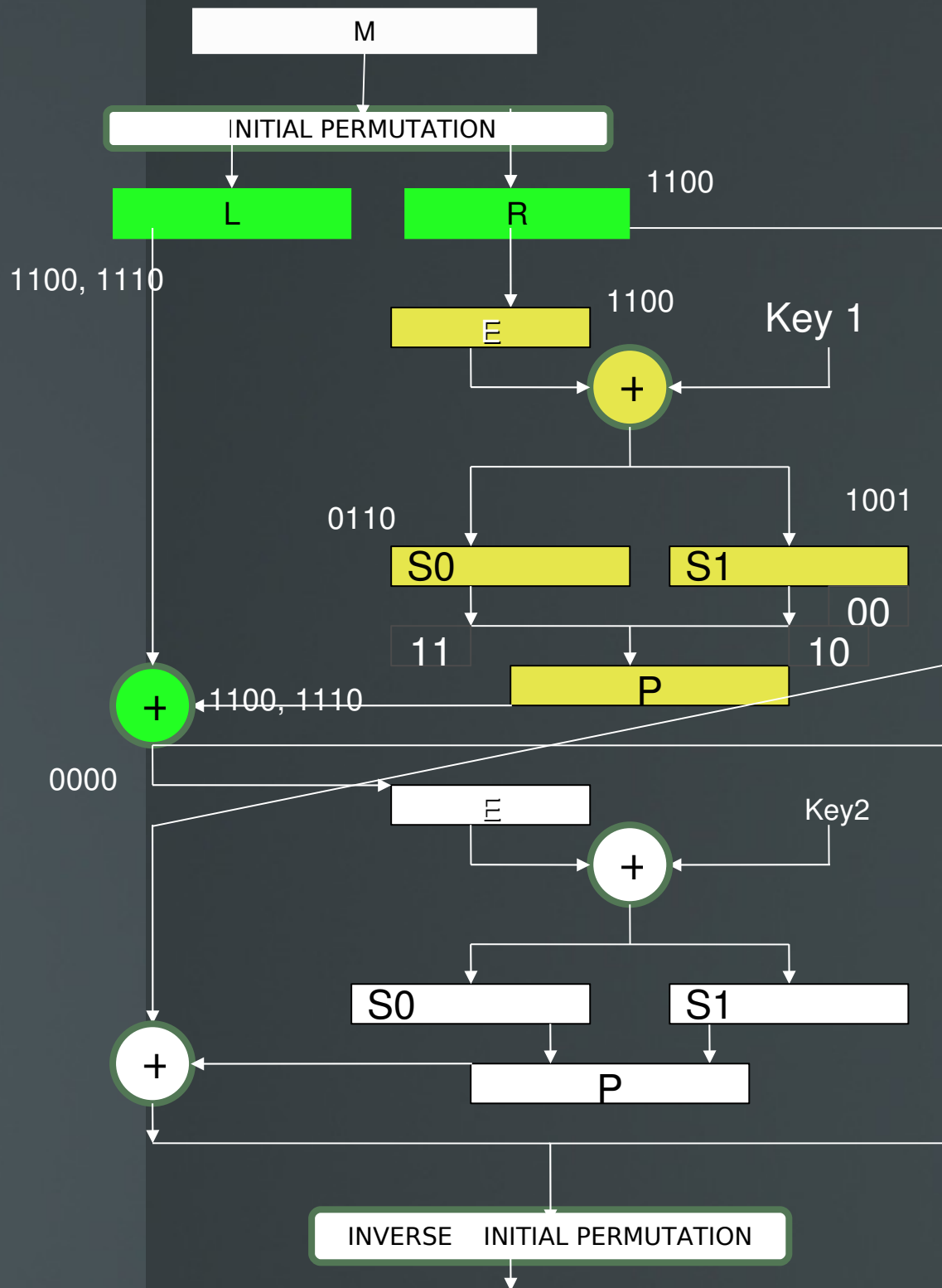
S1

<i>DX/DY</i>	0	1	2	3
0	16	0	0	0
1	2	8	2	4
2	0	6	4	6
3	4	2	8	2
4	2	0	10	4
5	2	4	2	8
6	0	10	0	6
7	8	2	4	2
8	4	6	0	6
9	8	2	4	2
10	2	0	10	4
11	0	6	4	6
12	0	6	4	6
13	6	0	6	4
14	10	4	2	0
15	0	8	2	4

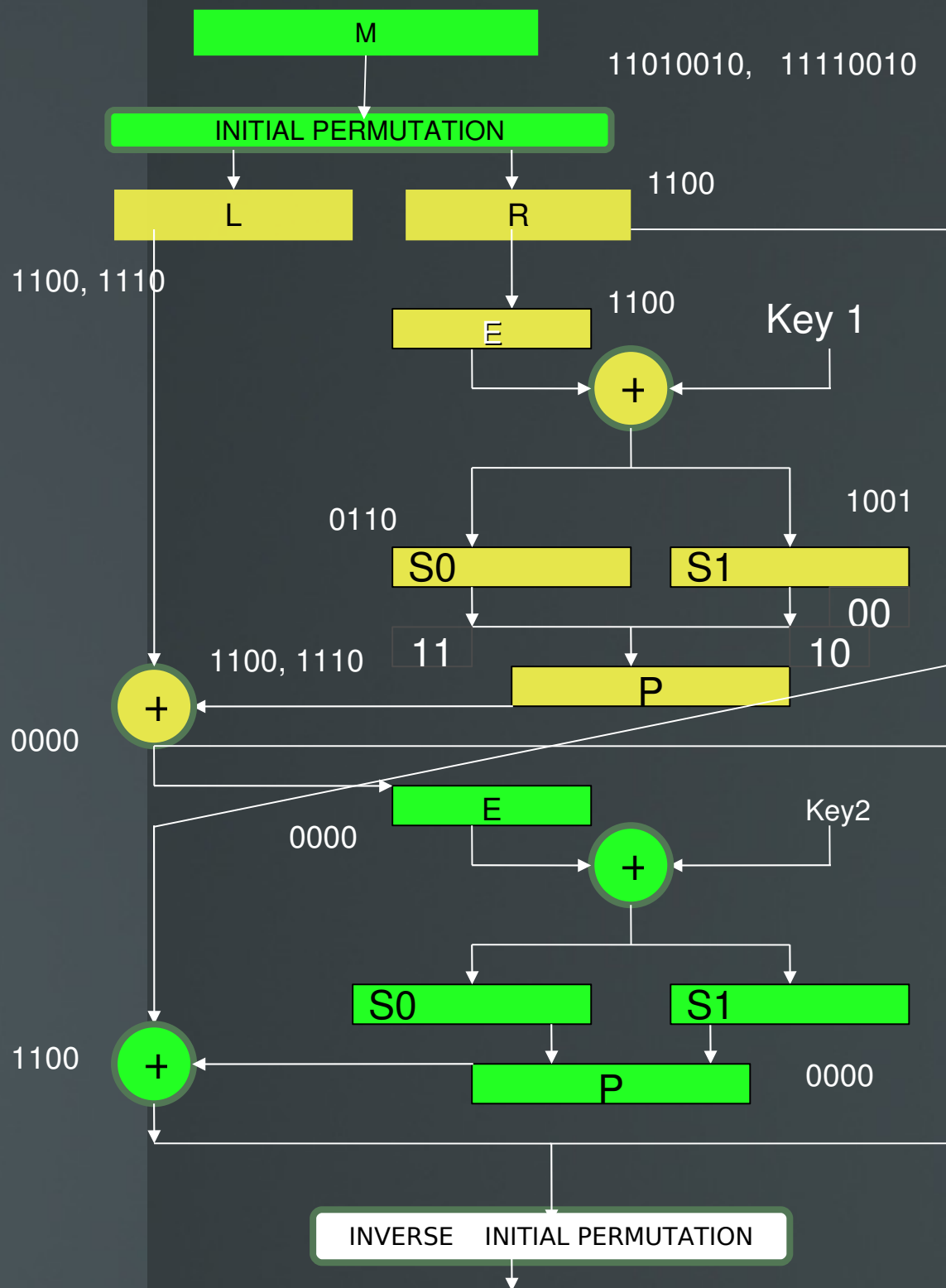
Differential Cryptanalysis



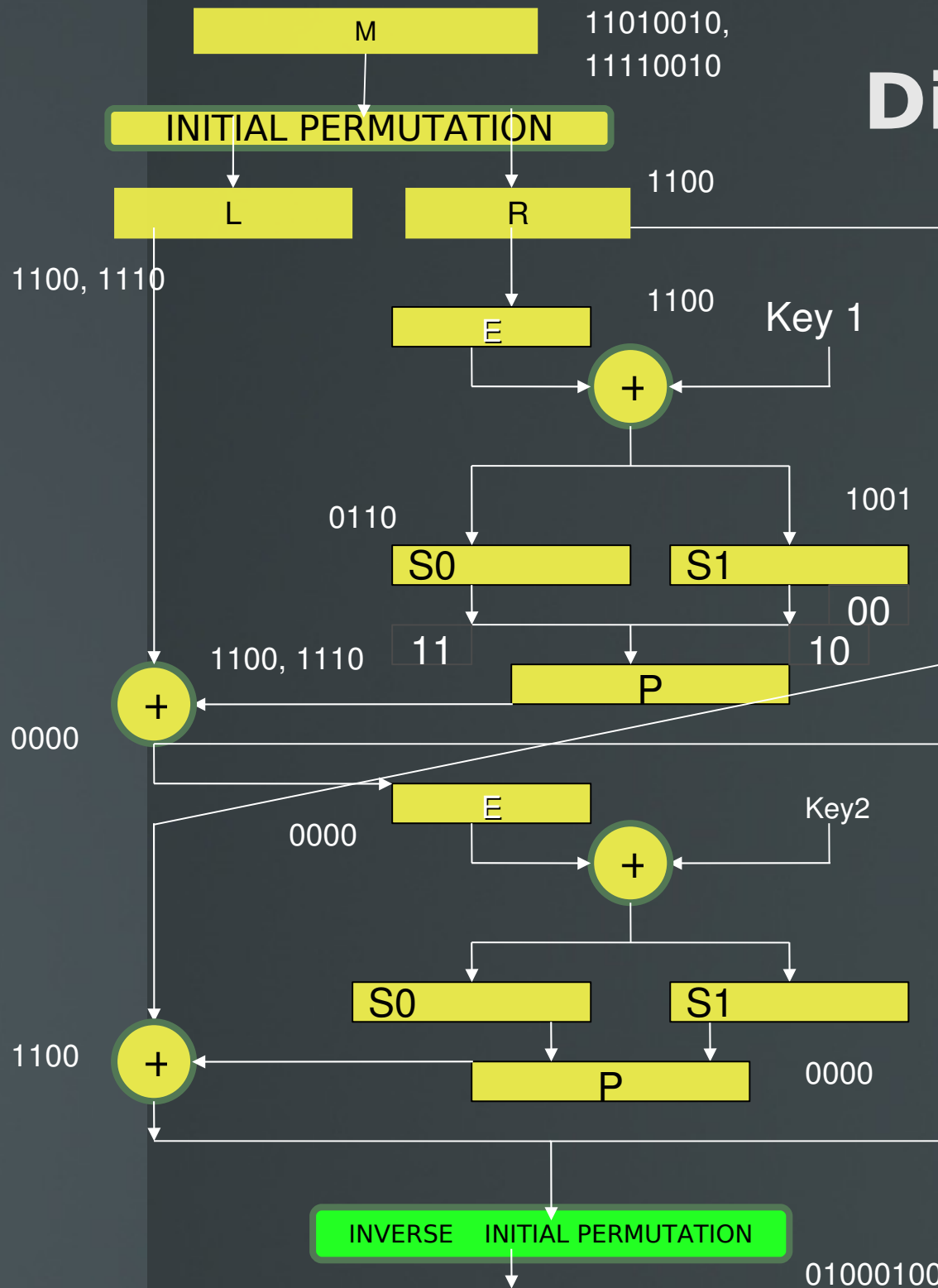
Differential Cryptanalysis



Differential Cryptanalysis



Diff. Cryptanalysis



del_C = 01000100

**del_P = 11010010,
11110010**

Calculate a Random P

$P_1 = P \oplus \text{del}_P1$

$P_2 = P \oplus \text{del}_P2$

$C = \text{Encrypt}(P, K)$

$C_1 = \text{Encrypt}(P_1, K)$

$C_2 = \text{Encrypt}(P_2, K)$

$\text{del_C1} = C_1 \oplus C$

$\text{del_C2} = C_2 \oplus C$

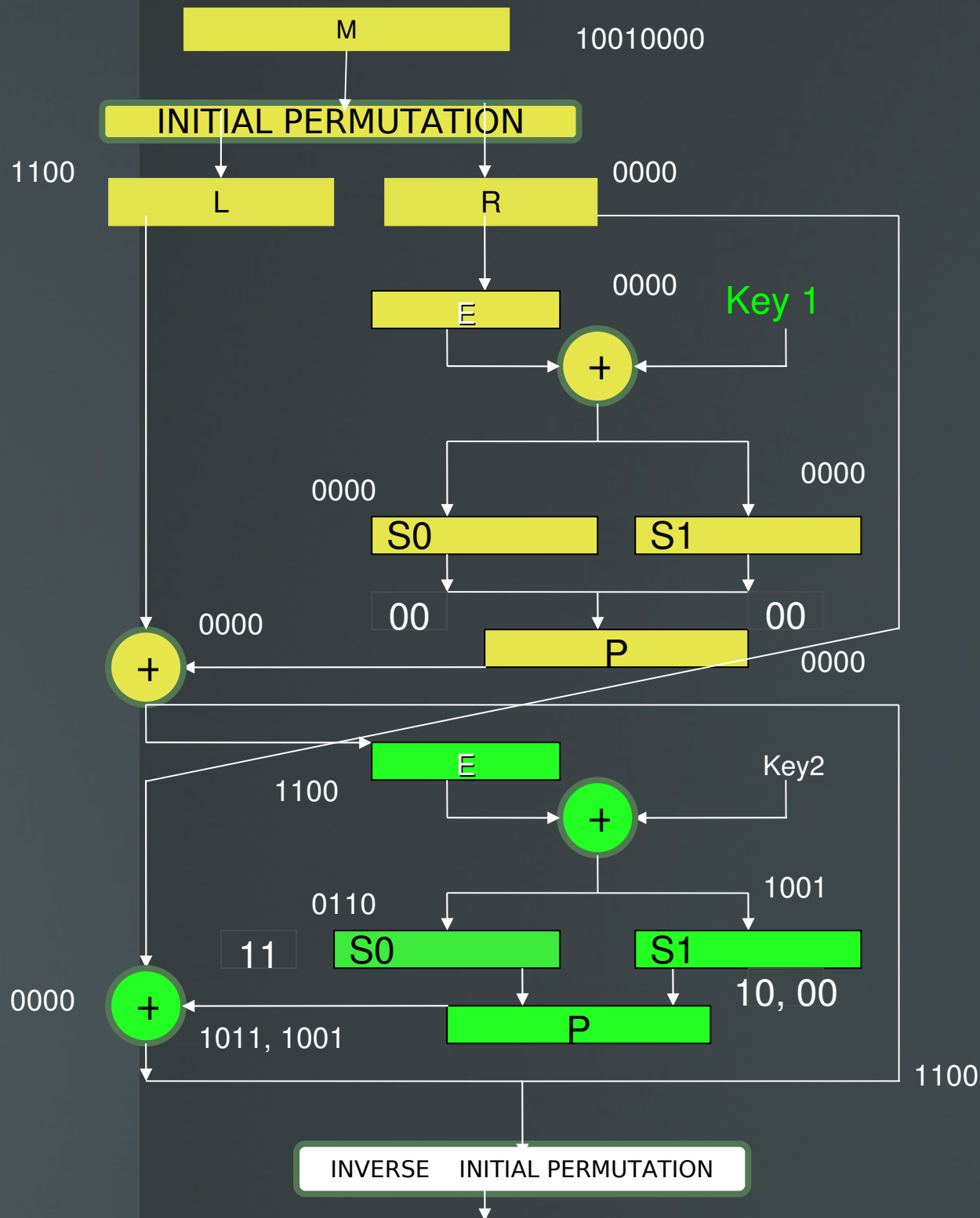
**If del_C1 or del_C2 is =
del_C characteristic has
occurred, if not repeat the
steps with different del_X
and del_Y**

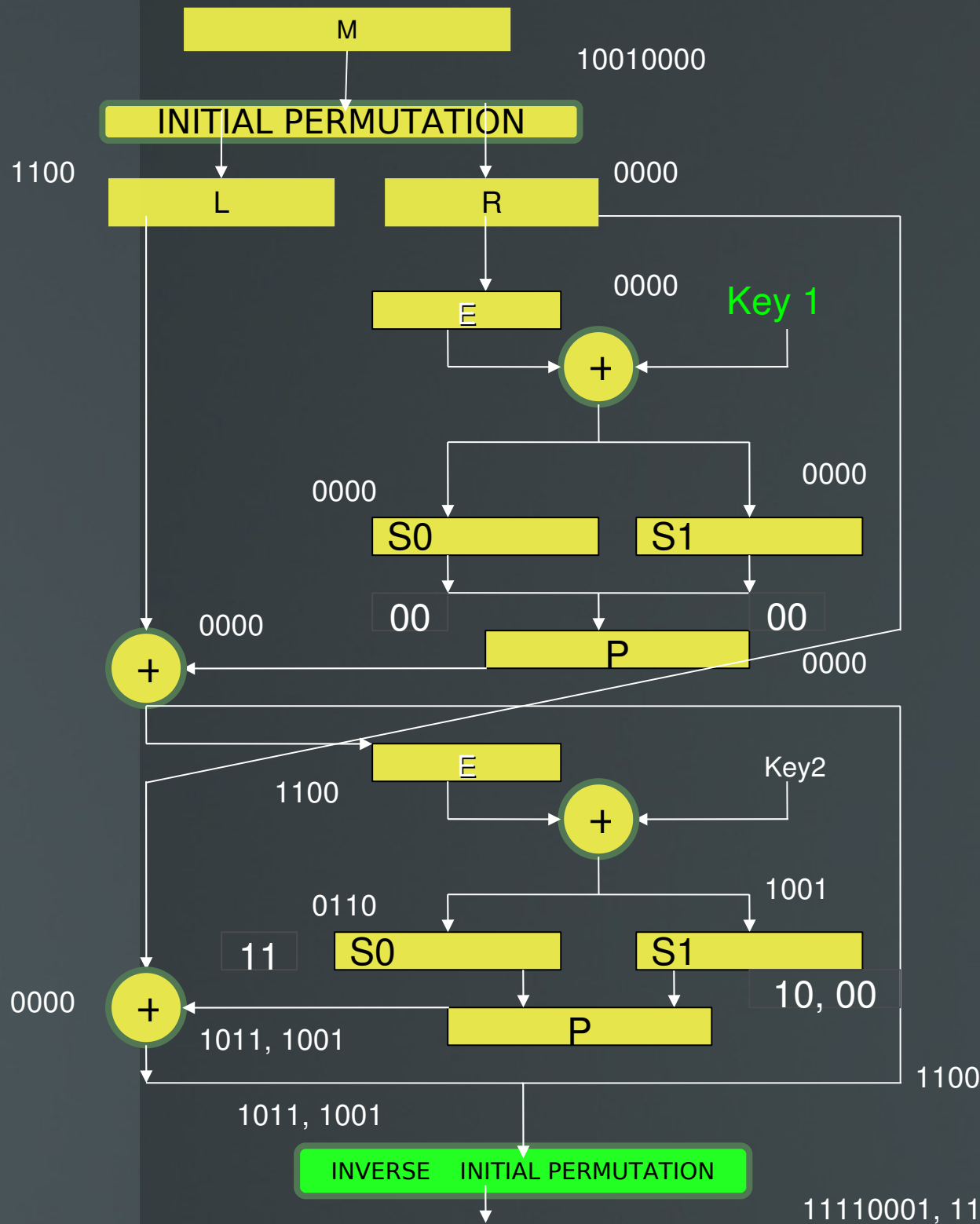
Determining the Key

Finding K1

- Select the Value of P that generated the differential characteristic
- Find the output of Expansion function
- Identify the set of X values(X_0 , X_1) that satisfy ΔE_X and ΔE_Y for both S_0 and S_1 .
- Left Half of K1 is $E_L \oplus X_0$, where E_L is the Left half of E, and X_0 is the input to S-box 0
- Right Half of K1 is $E_R \oplus X_1$, where E_R is the right half of E, and X_1 is the input to S-box 1.
- Iterate the above steps to find a Unique value of K1

Finding K2





$\text{del_C} = 11110001, 11010001$

$\text{del_P} = 10010000$

Calculate a Random P

$P_1 = P \oplus \text{del_P1}$

$C = \text{Encrypt}(P, K)$

$C_1 = \text{Encrypt}(P_1, K)$

$\text{del_C1} = C_1 \oplus C$

If $\text{del_C1} = \text{del_C}$
characteristic has occurred,
if not repeat the steps with
different del_X and del_Y

Determining the Key

Finding K2

- Use Key1 to calculate the output of the first round and find output of Expansion function for round 2
- Identify the set of X values(X_0 , X_1) that satisfy del_X and del_Y for both S_0 and S_1 .
- Left Half of K2 is $E_l \oplus X_0$, where E_l = Left half of E
- Right Half of K2 is $E_r \oplus X_1$, where E_r = Right half of E
- Iterate the above steps to find a Unique value of K2

Key Captured...!!!

Key1 = $K1_1, K1_2, K1_3, K1_4, K1_5, K1_6, K1_7, K1_8$

Key2 = $K2_1, K2_2, K2_3, K2_4, K2_5, K2_6, K2_7, K2_8$

Key is = $K1_1 K2_6 K1_6 K1_4 K2_4 K1_8 K1_2 K1_5 K1_3 K1_7$

Conclusion

This project demonstrates the successful use of differential cryptanalysis technique on Simplified DES.

Questions...!!