

Homework 3 of CS291A Introduction to Differential Privacy (Fall 2021)

University of California, Santa Barbara

Assigned on Nov 21, 2021(Sunday)

Due at 11:59 pm on Dec 3, 2021 (Wednesday)

Notes:

- Be sure to read “Policy on Academic Integrity” on the course syllabus.
 - There are *[100 points]* in this homework, and a bonus *[5 points]*.
 - You need to submit your homework via Gradescope.
 - Contact the instructor if you spot typos. Any updates or correction will be posted on the course Announcements page and piazza, so check there occasionally.
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0. Acknowledgment *[0 points]*

For each question in this HW, please list all your collaborators and reference materials (beyond those specified on the website) that were used for this homework.

1. **List of Collaborators** List the names of all people you have collaborated with and for which question(s).
2. **List of Acknowledgements.** If you find an assignment’s answer or use a another source for help, acknowledge for which question and provide an appropriate citation (there is no penalty, provided you include the acknowledgement). If not, then write “none”.

1 General facts about Differential Privacy [33 pts]

Please answer True or False and provide a short (one sentence explanation). (3 pts each)

- (a) Differential privacy prevents all harms that could incur to an individual when a dataset including this individual is analyzed (differentially privately).
- (b) Differentially Private machine learning algorithm cannot possibly predict my attribute accurately because otherwise it violates the DP guarantee.
- (c) Differential Privacy prevents attackers from identifying Bob, even if the rest of the dataset (except Bob) is known.

- (d) If one can implement homomorphic encryption computationally efficiently to train a machine learning model, then it provides stronger privacy guarantee than differential privacy.
- (e) DP algorithms must add noise.
- (f) Deterministic algorithms cannot be differentially private.
- (g) (ϵ, δ) -Approximate differential privacy provides DP guarantee with probability $1 - \delta$.
- (h) k -anonymity could fail completely upon composition, whereas differential privacy degrades more gracefully under composition.
- (i) It is often worthwhile to exploit the sparsity in the problem of differentially private histogram release. Adding noise only to those elements of the histogram that are non-zero helps to improve the utility of the DP release.
- (j) Gaussian mechanism is the only mechanism satisfying a linear upper bound of the Renyi differential privacy (i.e., Concentrated Differential Privacy).
- (k) Objective perturbation mechanism dominates output perturbation mechanism in differentially private (convex) empirical risk minimization.

2 Differentially private Linear Regression [21 pts]

In this question, you will see the connection of various mechanisms for differentially private linear regression and how they are related to each other. The goal is to solve

$$\min_{\theta} \sum_{i=1}^n (x_i^T \theta - y_i)^2 + \lambda \|\theta\|^2 = \|X\theta - y\|^2 + \lambda \|\theta\|^2$$

This this question, we do not concern choosing the parameter of the randomization to achieve a particular (ϵ, δ) -DP. You should just write down the form of the output using the native parameters of the algorithms.

- (a) (7 pts) Write down the output perturbation mechanism that adds noise to the solution in terms of X, y, λ , you may assume the noise you add is $\mathcal{N}(0, \sigma^2 I_d)$
- (b) (7 pts) Express the objective perturbation algorithm as a data-dependent noise-adding procedure. You may assume the noise vector $b \sim \mathcal{N}(0, \sigma^2 I_d)$
- (c) (7 pts) Express the posterior sampling algorithm as a data-dependent noise adding procedure. You may just express things as scale parameter γ such that we are outputting

$$\hat{\theta} \sim P(\theta|X, y) \propto \exp(-\gamma(\sum_{i=1}^n (x_i^T \theta - y_i)^2 + \lambda \|\theta\|^2)).$$

(Hint: the above samples from a multivariate Gaussian mechanism. What is the mean and what is the covariance matrix?)

3 Private selection and data-dependent DP [25 pts +5 bonus]

Consider the problem of private voting. Each voter can choose only one out of m candidates. The voting scores can be represented as a histogram. Assume that on our particular dataset, the most popular candidate attracts k more voters than the second most popular candidate.

We measure the utility by the probability of outputting the correct argmax.

- (a) Write down how Dist2Instability mechanism works and replicate its privacy analysis. (The cleanest description of it is in the recap of Lecture 16)
- (b) What is the utility of Dist2Instability mechanism with DP parameter (ϵ, δ) ? (Hint: Discuss what happens when you vary $0 \leq k \leq n$).
- (c) Consider Laplace mechanism (with DP parameter ϵ) for releasing the histogram, then output argmax as a post-processing. Lower bound the utility of this approach as a function of k ? (Notice that this mechanism is the same as report-noisy-max).
- (d) Consider exponential mechanism with DP parameter ϵ . Lower bound the utility of this approach as a function of k . (Hint: you could directly apply the theorem in Slide 12 of Lecture 5)
- (e) Draw the utility as a function of k (by hand or by matplotlib). When shall we use which algorithm when we need (ϵ, δ) -DP and want to have the largest utility (assuming we have a good guess what k is)?
- (f) (Bonus 5 pts) Derive the Gaussian mechanism version of Report-Noisy-Max and its utility.

4 Privacy amplification by sampling [21 pts]

Consider a dataset $x \in \{0, 1\}^n$ (represented as a binary bit vector), and the following algorithm:

1. Randomly select one coordinate of this dataset.
2. Run randomized response on the selected coordinate with parameter ϵ (i.e., output the correct value with probability $e^\epsilon/(1 + e^\epsilon)$)

In this question, we will work out the privacy parameter of this algorithm under the “Replace-One” neighboring relationship.

- (a) (7 pts) We will describe a dataset x under the assumption that the ordering of the coordinates does not matter. Let $0 \leq m \leq n$ be the number of 1s. Without loss of generality, assume the last coordinate is what differs between x and its neighbor x' . What is the probability of outputting 1?
- (b) (7 pts) Show that x and x' are ϵ' indistinguishable. Parameterize ϵ' by n, m, ϵ .
- (c) (7 pts) Show that this algorithm is ϵ'' -DP. Work out the parameter ϵ'' . Notice that you cannot use the general theorem of privacy amplification by sampling, but you may use it to check if your solution is correct.