

Lecture 14 Data-Dependent DP Algorithm design

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COMPUTER SCIENCE

UC SANTA BARBARA

Computing. ReInvented.

Recap: Differentially Private Machine Learning

- Private learning from a finite class is easy

$$\frac{\log(1/\delta)}{n \cdot \epsilon}$$

- Private learning from an infinite class is hard (in general)

Not privately learnable

- Let's restrict our attention to the Lipschitz losses

Recap: Convex empirical risk minimization

- Posterior sampling (i.e., exponential mechanism)
- Output perturbation / Objective perturbation
- NoisyGD and NoisySGD

(Gaussian Mechanism)^T

(Subsampled Gaussian)^T

$f(b, \theta)$

Recap: NoisyGD summary

	Lipschitz + convex	Lipschitz + Smooth + convex	Smooth + Lipschitz + convex + GLM
Output Pert	$\frac{d^{1/4}L\ \theta^*\ \log(\frac{1}{\delta})^{1/4}}{n^{1/2}\epsilon^{1/2}}$	$\frac{d^{1/3}\beta^{1/3}L^{2/3}\ \theta^*\ ^{4/3}\log(\frac{1}{\delta})^{1/3}}{n^{2/3}\epsilon^{2/3}}$	Same as left
ObjPert	Not applicable	$\frac{dL\ \theta^*\ \sqrt{\log(\frac{1}{\delta})}}{n\epsilon}$ Lower order terms and dependence on β hidden.	$\frac{\sqrt{d}L\ \theta^*\ \sqrt{\log(\frac{1}{\delta})}}{n\epsilon}$
NoisyGD	$\frac{\sqrt{d}L\ \theta^*\ \sqrt{\log(\frac{1}{\delta})}}{n\epsilon}$	$\frac{\sqrt{d}L\ \theta^*\ \sqrt{\log(\frac{1}{\delta})}}{n\epsilon}$	$\frac{\sqrt{d}L\ \theta^*\ \sqrt{\log(\frac{1}{\delta})}}{n\epsilon}$

	Lipschitz + Strongly convex	Lipschitz + Smooth + Nonconvex
NoisyGD	$\frac{dL^2\log(1/\delta)}{n\lambda\epsilon^2}$	$\frac{\sqrt{n\beta d}L^2(f(\theta_1) - f^*)\log(1/\delta)}{n\epsilon}$ Stationary point convergence

Recap: Comparing NoisyGD and NoisySGD computationally

- Both optimal information-theoretically.
 - If we ignore computation and add very large noise, but use infinitesimal step-size
- Table to compare computation
 - in terms of **the number of incremental gradient calls** to achieve **information theoretic limit** up to a constant

	Lipschitz + Smooth + Convex	Lipschitz + Convex	Lipschitz + Strongly convex
NoisyGD	$\frac{n^2 \beta \ x_1 - x^*\ \sqrt{\rho}}{\sqrt{d} L}$	$\frac{n^3 \rho}{d}$	$\frac{n^3 \rho}{\lambda}$
NoisySGD	$\frac{n^{3/2} \beta^{1/2} \ x_1 - x^*\ \rho^{1/2}}{d^{1/4} L^{1/2}} + \frac{n^2 \rho}{d}$	$\frac{n^2 \rho^{3/4}}{d^{1/2}} + \frac{n^2 \rho}{d}$	$\frac{n^2 \rho^{3/4}}{d^{1/2}} + \frac{n^2 \rho}{d}$

Open problem: what is the optimal computational complexity?

Recap: Deep Learning with DP

- NoisySGD with per-example gradient clipping
 - The only practical / popular algorithm
 - Empirical research questions: What are tricks to improve NoisySGD?
 - Theoretical open problem: what exactly is the effect of gradient clipping in training? How does it work?
- Assume access to some (unlabeled) public data
 - Private Aggregation of Teacher Ensembles.
 - PrivateKNN

Very few public data points are needed in PATE... also it learns all VC-classes in the realizable setting.

Table 1: Summary of our results: excess risk bounds for PATE algorithms.

Algorithm	PATE (Gaussian Mech.)	PATE (SVT-based)		PATE (Active Learning)
	Papernot et al. [2017]	Bassily et al. [2018a]	This paper	This paper
Realizable	$\tilde{O}\left(\frac{d}{(n\epsilon)^{2/3}} \vee \frac{d}{m}\right) = \mathcal{O}$	$\tilde{O}\left(\frac{d}{(n\epsilon)^{2/3}} \vee \sqrt{\frac{d}{m}}\right)$	$\tilde{O}\left(\frac{d^{3/2}}{n\epsilon} \vee \frac{d}{m}\right)$	$\tilde{O}\left(\frac{d^{3/2}\theta^{1/2}}{n\epsilon} \vee \frac{d}{m}\right)$
τ -TNC	$\tilde{O}\left(\left(\frac{d^{3/2}}{n\epsilon}\right)^{\frac{2\tau}{4-\tau}} \vee \frac{d}{m}\right)$	same as agnostic	$\tilde{O}\left(\left(\frac{d^{3/2}}{n\epsilon}\right)^{\frac{\tau}{2-\tau}} \vee \frac{d}{m}\right)$	$\tilde{O}\left(\left(\frac{d^{3/2}\theta^{1/2}}{n\epsilon}\right)^{\frac{\tau}{2-\tau}} \vee \frac{d}{m}\right)$
Agnostic (vs h^*)	$\Omega(\text{Err}(h^*))$ required.	$13\text{Err}(h^*) + \tilde{O}\left(\frac{d^{3/5}}{n^{2/5}\epsilon^{2/5}} \vee \sqrt{\frac{d}{m}}\right)$	$\Omega(\text{Err}(h^*))$ required.	$\Omega(\text{Err}(h^*))$ required.
Agnostic (vs h_∞^{agg})	-	-	Consistent under weaker conditions.	-

This lecture

- Going beyond the worst case!
- Smoothed Sensitivity and the Median
- Concentrated DP of Smoothed Sensitivity-based algorithm

Reading materials

- Nissim, Raskhodnikova, Smith 2011 “Smooth Sensitivity and Sampling in Privacy Data Analysis”:
<https://cs-people.bu.edu/ads22/pubs/NRS07/NRS07-full-draft-v1.pdf>
- Bun and Steinke 2019: “Average cases averages”
<https://arxiv.org/abs/1906.02830>

Recap: Private query release

- For example, linear queries

$$\underline{f(x) = \frac{1}{n} \sum_{i=1}^n \phi(x_i)}$$

$$X = \{x_1, \dots, x_n\}$$

n is public, replace one x_i with x'_i
 $X' = \{x_1, \dots, x'_i, \dots, x_n\}$
 $d(X, X') = \frac{1}{n}$

- Laplace mechanism / Gaussian mechanism

$$A(x) = f(x) + \text{Lap}\left(\frac{\Delta_1(f)}{\epsilon}\right) \leftarrow \epsilon\text{-DP}$$

$$\text{Gaussm}(0, G^2) \leftarrow \frac{\Delta_2^2(f)}{2G^2} \text{-CDP}$$

- Global sensitivity

$$\Delta_1(f) \Rightarrow \max_{d(x, x') \leq 1} \|f(x) - f(x')\|_1$$

$$\Delta_2(f) \Rightarrow \|f(x) - f(x')\|_2$$

An example when the global sensitivity approach is very inefficient

- Median query: $X := \text{Dataset} = \{x_1, \dots, x_n\}$ $x_i \in [0, 1]$
 Assume n is an odd number
 $x_{(1)}, x_{(2)}, \dots, x_{(\frac{n+1}{2})}, \dots, x_{(n)}$ order statistics
 $f_{\text{med}}(X) = x_{(\frac{n+1}{2})}$
 Sort
 Median

- Example:
 $\{0, 0, 0, \dots, 0, 1, 1, \dots, 1\}$
 Median
 $\{0, 0, 0, \dots, 1, 1, 1, \dots, 1\}$
 Median

$$|f_{\text{med}}(x) - f_{\text{med}}(x')| = 1$$

Another example: linear regression

$$\min_{\theta} \frac{1}{m} \sum_{i=1}^m (x_i^T \theta - y_i)^2 = \|X\theta - \vec{y}\|_2^2 + \lambda \|\theta\|_2^2$$

- The output perturbation mechanism, revisited

$$(X, \vec{y}) \quad \theta^* = (X^T X)^{-1} X^T \vec{y}$$

Δ_{CS} is unbounded

$$\sum x_i (x_i^T \theta, y_i)$$

$$((X, x), (\vec{Y}, y)) \quad \theta^* = (X^T X + \lambda I)^{-1} (X^T \vec{Y} + \lambda y)$$

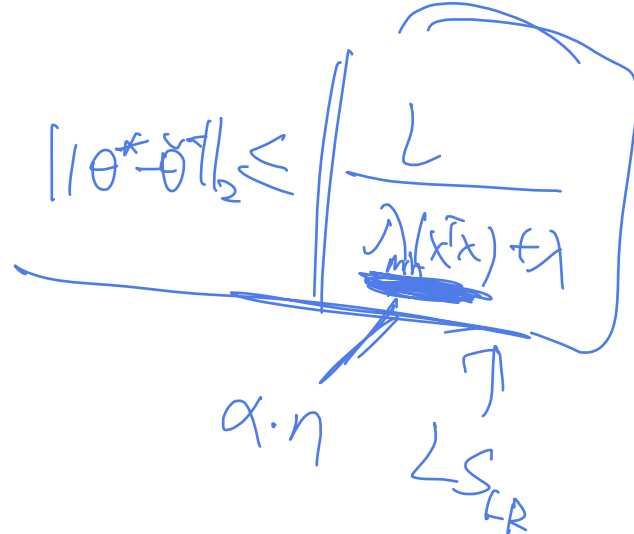
$$\|\theta^* - \theta^*\|_2 \leq \frac{\|X\|_F \|\theta\|_2 + \|X\|_F \|y\|_2}{\lambda}$$

- The global sensitivity approach does not exploit the fact that the input dataset is well-conditioned

$$\text{if } X^T X \succeq \alpha \cdot n \cdot I$$

\Rightarrow

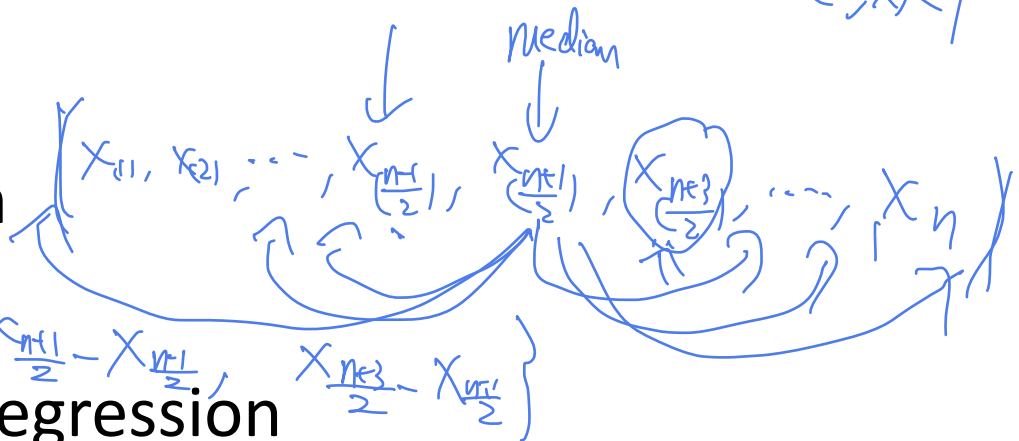
$$\|\theta^* - \theta^*\|_2 \leq$$



Local Sensitivity measures the stability of a query at a particular input dataset.

$$LS_q(x) = \max_{x'} \left\{ |q(x) - q(x')| : \underbrace{x' \sim x}_{d(x, x') \leq 1} \right\}.$$

- Example: median



$$LS_{f_{\text{med}}}(x) = \max \left\{ x_{(n/2+1)} - x_{(n/2)}, x_{(n/2+2)} - x_{(n/2+1)} \right\}$$

- Example: linear regression

$$\theta^* = (X^T X)^{-1} X^T y$$

$$\frac{\partial L}{\partial \text{dim}(X^T X)}$$

The issue of calibrating noise to local sensitivity

- Example of the median

$$x = (0, 0, \dots, 0, \overset{\text{noise}}{0}, 0, \dots, 1) \quad \text{Median}$$

$$x' = (0, 0, \dots, 0, 0, \overset{0.20}{0}, \dots, 1) \quad \text{Median}$$

$$\text{Lap}\left(\frac{LS}{\epsilon}\right)$$

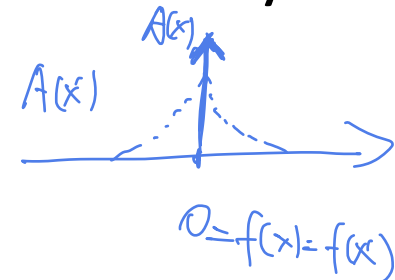
$$LS_{f_{\text{med}}}(x) = 0$$

$$f_{\text{ma}}(x) = 0$$

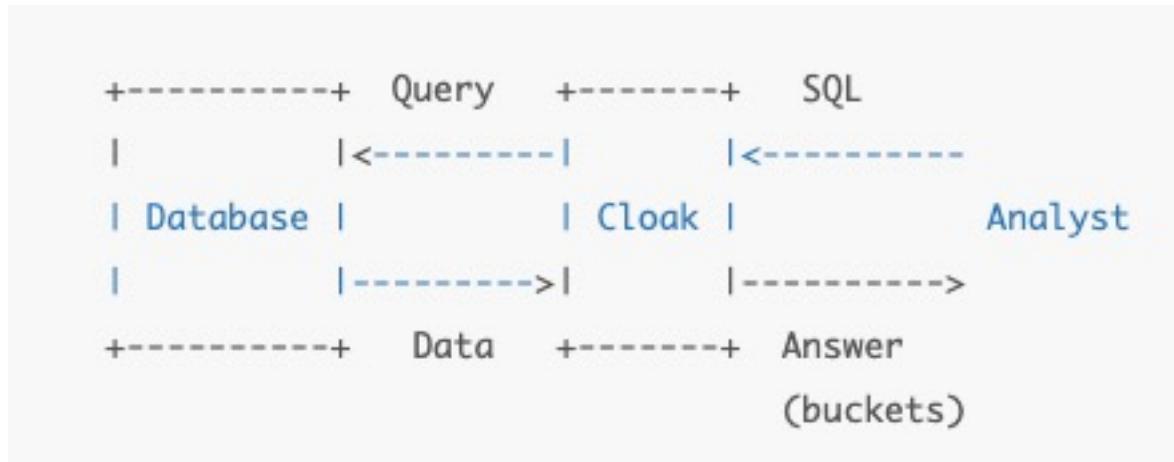
$$f_{\text{ma}}(x) = 0$$

$$LS_{f_{\text{med}}}(x') = \text{? } 0.20$$

- In conclusion: the magnitude of the noise may reveal sensitive information!



Diffix and “Sticky Noise”



*Implementing a bunch of heuristics to protect against known attacks.
Decide how much noise to add by the specific dataset and how sensitive the question is.*

From Creator of Diffix:

“anonymizing SQL interface [that] sits in front of your data and enables you to conduct ad hoc analytics — fully privacy preserving and GDPR-compliant.”

Attack on Diffix

When the Signal is in the Noise: Exploiting Diffix's Sticky Noise

Andrea Gadotti^{*a}, Florimond Houssiau^{*a}, Luc Rocher^{*a,b}, Benjamin Livshits^a, and Yves-Alexandre de Montjoye^{†a}

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Link to the paper:

https://www.usenix.org/system/files/sec19fall_gadotti_prepub.pdf

Also see this nice post: <https://differentialprivacy.org/diffix-attack/>

“Data-dependent DP mechanism” aims at **more stably** calibrating noise to local sensitivity (at least for query releases)

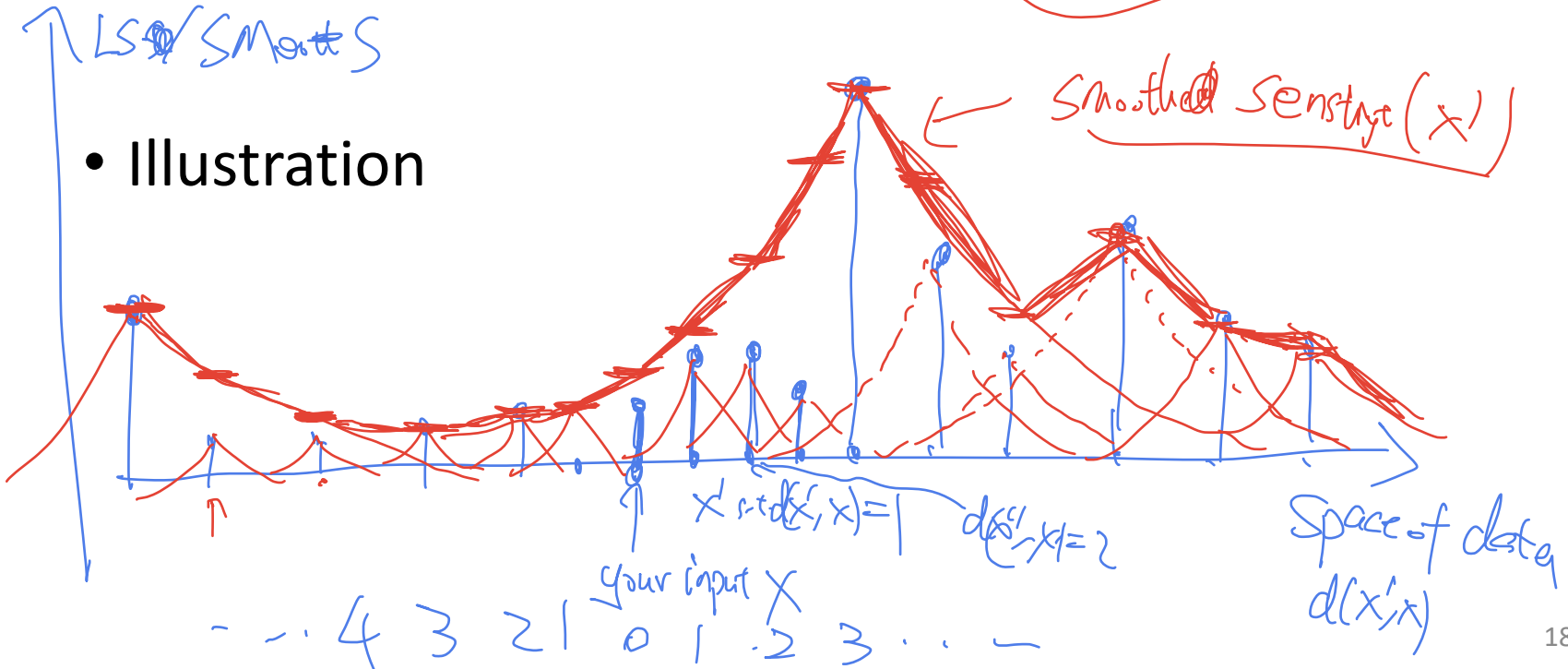
- A number of different approaches:
 - Smooth sensitivity
 - Propose-test-release
 - Privately bounding the local-sensitivity
 - Stability-based query release (Distance2Stability)

Smooth Sensitivity

$$S^*(y) \stackrel{\text{def}}{=} (LS_f(x') \cdot e^{-\beta d(x',y)})$$

DEFINITION 2.2 (SMOOTH SENSITIVITY). For $\beta > 0$, the β -smooth sensitivity of f is

$$S_{f,\beta}^*(x) = \max_{y \in D^n} \left(LS_f(y) \cdot e^{-\beta d(x,y)} \right).$$



Smooth sensitivity satisfies a **smoothing property**, and it is the **optimal bound** satisfying this property

- Two properties that one should satisfy to smooth out the local sensitivity

$$\forall x \in D^n : \quad \underline{S(x) \geq LS_f(x)} ; \quad \checkmark$$

$$\forall x, y \in D^n, d(x, y) = 1 : \quad \underline{S(x) \leq e^\beta S(y)} \quad \left[\beta \text{ Smooth property} \right]$$

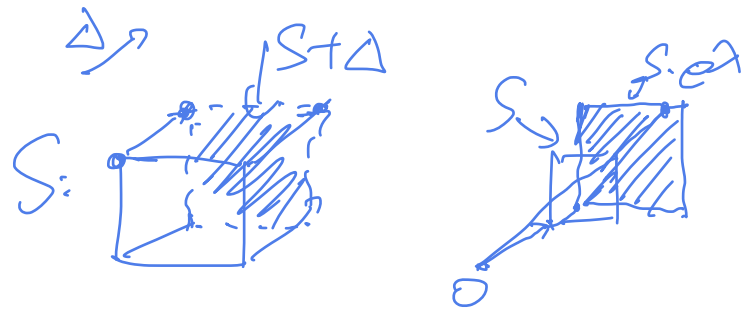
- Smooth sensitivity is the optimal bound

LEMMA 2.3. $S_{f,\beta}^*$ is a β -smooth upper bound on LS_f . In addition, $S_{f,\beta}^*(x) \leq S(x)$ for all $x \in D^n$ for every β -smooth upper bound S on LS_f .

$d(x,y)=1, x' \text{ s.t. } S_{f,\beta}^*(x) = LS_f(x') e^{-d(x,x')}
$$S^*(y) \geq L(x') e^{-\beta d(y,x)}$$

$$\geq \frac{L(x') \cdot e^{-\beta d(x,x') - \beta}}{e^{-\beta d(x,x')}} = e^{-\beta} L(x') \cdot e^{-\beta d(x,x')} = e^{-\beta} S^*(x) d$$$

What noise to add?



Notation. For a subset \mathcal{S} of \mathbb{R}^d , we write $\mathcal{S} + \Delta$ for the set $\{z + \Delta \mid z \in \mathcal{S}\}$, and $e^\lambda \cdot \mathcal{S}$ for the set $\{e^\lambda \cdot z \mid z \in \mathcal{S}\}$. We also write $a \pm b$ for the interval $[a - b, a + b]$.

Definition 2.5 (Admissible Noise Distribution). A probability distribution on \mathbb{R}^d , given by a density function h , is (α, β) -admissible (with respect to ℓ_1) if, for $\alpha = \alpha(\epsilon, \delta)$, $\beta = \beta(\epsilon, \delta)$, the following two conditions hold for all $\Delta \in \mathbb{R}^d$ and $\lambda \in \mathbb{R}$ satisfying $\|\Delta\|_1 \leq \alpha$ and $|\lambda| \leq \beta$, and for all measurable subsets $\mathcal{S} \subseteq \mathbb{R}^d$:

Sliding Property:
$$\Pr_{Z \sim h} [Z \in \mathcal{S}] \leq e^{\frac{\epsilon}{2}} \cdot \Pr_{Z \sim h} [Z \in \mathcal{S} + \Delta] + \frac{\delta}{2}.$$

Dilation Property:
$$\Pr_{Z \sim h} [Z \in \mathcal{S}] \leq e^{\frac{\epsilon}{2}} \cdot \Pr_{Z \sim h} [Z \in e^\lambda \cdot \mathcal{S}] + \frac{\delta}{2}.$$

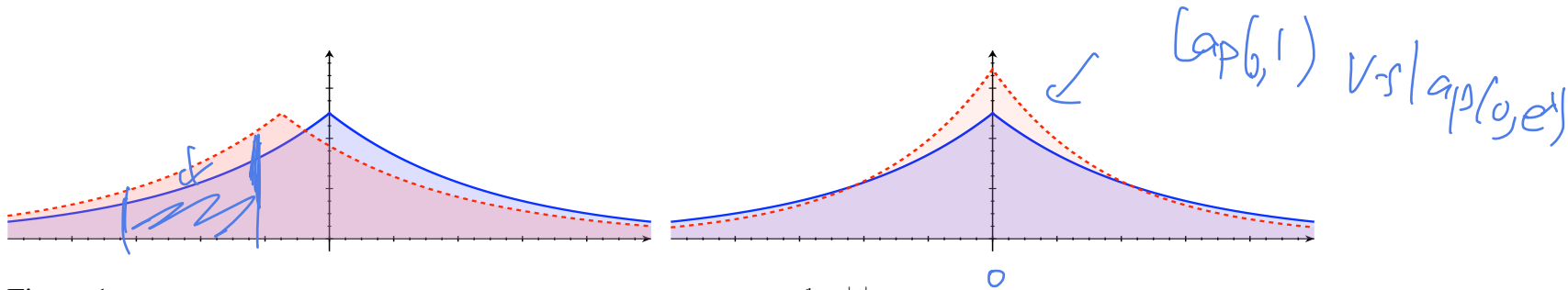


Figure 1: Sliding and dilation for the Laplace distribution with p.d.f. $h(z) = \frac{1}{2}e^{-|z|}$, plotted as a solid line. The dotted lines plot the densities $h(z + 0.3)$ (left) and $e^{0.3}h(e^{0.3}z)$ (right).

- Then $\mathcal{A}(x) = f(x) + \frac{S(x)}{\alpha} \cdot Z$ satisfies (ϵ, δ) -DP.

Privacy analysis

$$P(A(x) \in S) \leq e^\epsilon P(A(x') \in S) + \delta$$

$$P(A(x) \in S) = P\left(f(x) + \frac{S(x)}{\alpha} \cdot z \in S\right) = P\left(z \in \frac{\alpha(S - f(x))}{S(x)}\right)$$

$\frac{f(x) - f(x')}{S(x)}$
 $\frac{f(x) - f(x')}{S(x)}$

$$\leq e^{\frac{\alpha \Delta}{2\sigma}} P\left(z \in \frac{\alpha(S - f(x'))}{S(x)}\right) + \frac{\delta}{2}$$

$$\Delta = \frac{\alpha |f(x) - f(x')|}{S(x)}$$

$$\|\Delta\|_1 = \frac{\alpha \|f(x) - f(x')\|_1}{S(x)} \leq \frac{\alpha L S(x)}{S(x)} \leq \alpha$$

$$= e^{\frac{\alpha \Delta}{2\sigma}} \left(P\left(z \in \frac{S(x)}{S(x)} \cdot \frac{\alpha(S - f(x'))}{S(x)}\right) + \frac{\delta}{2} \right)$$

$$e^{\frac{\alpha \Delta}{2\sigma}} = \frac{S(x')}{S(x)} \leq e^\beta$$

$\lambda \leq \beta$

$$\leq e^{\frac{\alpha \Delta}{2\sigma}} \left(e^{\frac{\alpha \Delta}{2\sigma}} P\left(z \in \frac{\alpha(S - f(x'))}{S(x)}\right) + \frac{\delta}{2} \right) + \frac{\delta}{2}$$

$$= e^{\frac{\alpha \Delta}{2\sigma}} \left(e^{\frac{\alpha \Delta}{2\sigma}} P(f(x') \in S) + \frac{\delta}{2} \right) + \frac{\delta}{2}$$

$$\leq e^\epsilon P(f(x) \in S) + (e^{\frac{\alpha \Delta}{2\sigma}} + 1) \frac{\delta}{2}$$

Example: Cauchy-Noise, Laplace-noise, Gaussian noise

Lemma 2.7. For any $\gamma > 1$, the distribution with density $h(z) \propto \frac{1}{1+|z|^\gamma}$ is $(\frac{\epsilon}{2(\gamma+1)}, \frac{\epsilon}{2(\gamma+1)})$ -admissible (with $\delta = 0$). Moreover, the d -dimensional product of independent copies of h is $(\frac{\epsilon}{2(\gamma+1)}, \frac{\epsilon}{2d(\gamma+1)})$ -admissible.

Lemma 2.9. For $\epsilon, \delta \in (0, 1)$, the d -dimensional Laplace distribution, $h(z) = \frac{1}{2^d} \cdot e^{-\|z\|_1}$, is (α, β) -admissible with $\alpha = \frac{\epsilon}{2}$, and $\beta = \frac{\epsilon}{2\rho_{\delta/2}(\|Z\|_1)}$, where $Z \sim h$. In particular, it suffices to use $\alpha = \frac{\epsilon}{2}$ and $\beta = \frac{\epsilon}{4(d+\ln(2/\delta))}$. For $d = 1$, it suffices to use $\beta = \frac{\epsilon}{2\ln(2/\delta)}$.

Lemma 2.10 (Gaussian Distribution). For $\epsilon, \delta \in (0, 1)$, the d -dimensional Gaussian distribution, $h(z) = \frac{1}{(2\pi)^{d/2}} \cdot e^{-\frac{1}{2}\|z\|_2^2}$, is (α, β) -admissible for the Euclidean metric with $\alpha = \frac{\epsilon}{5\rho_{\delta/2}(Z_1)}$, and $\beta = \frac{\epsilon}{2\rho_{\delta/2}(\|Z\|_2^2)}$, where $Z = (Z_1, \dots, Z_d) \sim h$.

In particular, it suffices to take $\alpha = \frac{\epsilon}{5\sqrt{2\ln(2/\delta)}}$ and $\beta = \frac{\epsilon}{4(d+\ln(2/\delta))}$.

How to compute smooth sensitivity (or an upper bound of it?)

DEFINITION 2.2 (SMOOTH SENSITIVITY). For $\beta > 0$, the β -smooth sensitivity of f is

$$S_{f,\beta}^*(x) = \max_{y \in D^n} \left(LS_f(y) \cdot e^{-\beta d(x,y)} \right).$$

- An easier way to solve this optimization

$$S_{f,\epsilon}^*(x) = \max_{k=0,1,\dots,n} e^{-k\epsilon} \left(\max_{y: d(x,y)=k} LS_f(y) \right)$$

$$GS(f) \leq \bigwedge$$

$$\leq \max_{k=0,1,\dots,m} e^{-k\epsilon} \left(\max_{y: d(x,y)=k} LS_f(y) \right)$$

$m < n$

$$e^{-(m+1)\epsilon} \bigwedge$$

Example: smooth sensitivity of median

- Recall: $0 \leq x_1 \leq \dots \leq x_n \leq \Lambda.$

$$GS_{f_{med}} = \Lambda \quad LS_{f_{med}} = \max(x_m - x_{m-1}, x_{m+1} - x_m) \text{ for } m = \frac{n+1}{2}$$

- Now: $S_{f,\epsilon}^*(x) = \max_{k=0,1,\dots,n} e^{-k\epsilon} \left(\max_{y: d(x,y)=k} LS_f(y) \right)$

$$\max_{y: d(x,y) \leq k} LS(y) = \max_{0 \leq t \leq k+1} (x_{m+t} - x_{m+t-k-1}).$$

$O(n^2)$
 \Downarrow
 D.P.
 $O(n \log n)$



Checkpoint: smooth sensitivity

- We cannot calibrate noise to local sensitivity
 - Because noise-level itself may be sensitive
- Idea: construct smooth upper bound of local sensitivity
- Noise that satisfies stability under “translation” and “scaling” are admissible

Concentrated DP analysis of Smoothed Sensitivity

- Adding log-normal noise

$$Z = X \cdot e^{\sigma Y}$$

$$\text{Var}(Z) = \frac{1}{3} e^{2\sigma^2}$$

- X drawn from Laplace and Y from a standard Normal.

Proposition 3. Let $f : \mathcal{X}^n \rightarrow \mathbb{R}$ and let $Z \leftarrow \text{LLN}(\sigma)$ for some $\sigma > 0$. Then, for all $s, t > 0$, the algorithm $M(x) = f(x) + \frac{1}{s} \cdot S_f^t(x) \cdot Z$ guarantees $\frac{1}{2}\varepsilon^2$ -CDP for $\varepsilon = t/\sigma + e^{3\sigma^2/2}s$.

Summary of the noises that are known to work

- Cauchy distribution
 - Student t-distribution
- } → ϵ -DP
- Laplace-log-normal
 - Uniform-log-normal
 - Arcsinh-normal
- | → ρ -CDP
- Gaussian
 - Laplace
- | → (ϵ, δ) -DP

Sketch of the proof for the Laplace-Log-Normal

- Let's say for all neighboring datasets

$$|f(x) - f(x')| \leq g(x) \quad \text{and} \quad e^{-t}g(x) \leq g(x') \leq e^t g(x).$$

- **Algorithm:** $M(x) = f(x) + \frac{g(x)}{s} \cdot Z$ for $Z \leftarrow \text{LLN}(\sigma)$.

- **We have that** $D_\alpha(M(x) \| M(x')) = D_\alpha \left(Z \left\| \frac{f(x') - f(x)}{g(x)} \cdot s + \frac{g(x')}{g(x)} \cdot Z \right. \right)$.

Technical tools

- Group privacy for CDP:

Lemma 11. *Let P, Q, R be probability distributions. Suppose $D_\alpha(P\|R) \leq a \cdot \alpha$ and $D_\alpha(R\|Q) \leq b \cdot \alpha$ for all $\alpha \in (1, \infty)$. Then, for all $\alpha \in (1, \infty)$,*

$$D_\alpha(P\|Q) \leq \alpha \cdot (\sqrt{a} + \sqrt{b})^2 \leq 2\alpha \cdot (a + b).$$

- Decompose what we want to bound

$$D_\alpha(Z\|e^t Z + s)$$

$$D_\alpha(e^t Z + s\|Z)$$

Bounding the two parts separately

Lemma 19. *Let $Z \leftarrow \text{LLN}(\sigma)$ for $\sigma > 0$. Let $t \in \mathbb{R}$ and $\alpha \in (1, \infty)$. Then*

$$D_\alpha(Z \| e^t Z) \leq \frac{\alpha t^2}{2\sigma^2}.$$

- **Proof:**

$$D_\alpha(Z \| e^t Z) = D_\alpha(Xe^{\sigma Y} \| Xe^{\sigma Y+t}) \leq \sup_x D_\alpha(xe^{\sigma Y} \| xe^{\sigma Y+t}) \leq D_\alpha(\sigma Y \| \sigma Y + t).$$

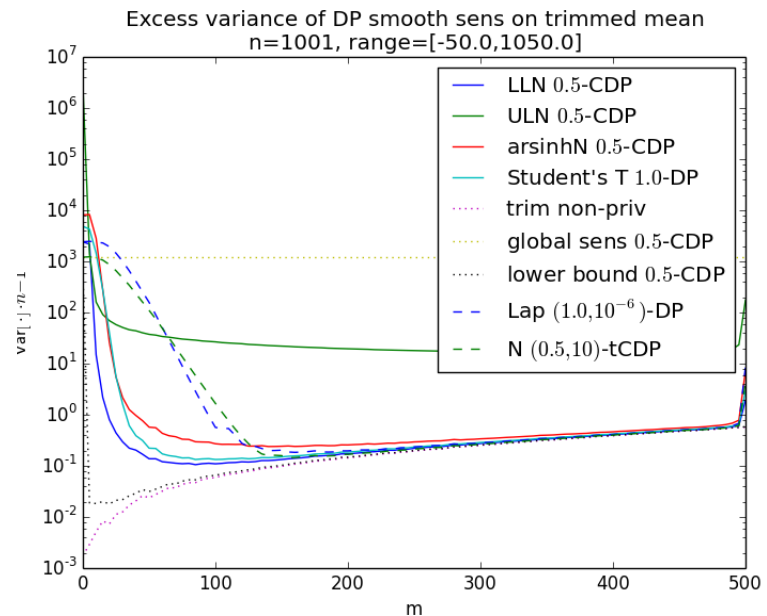
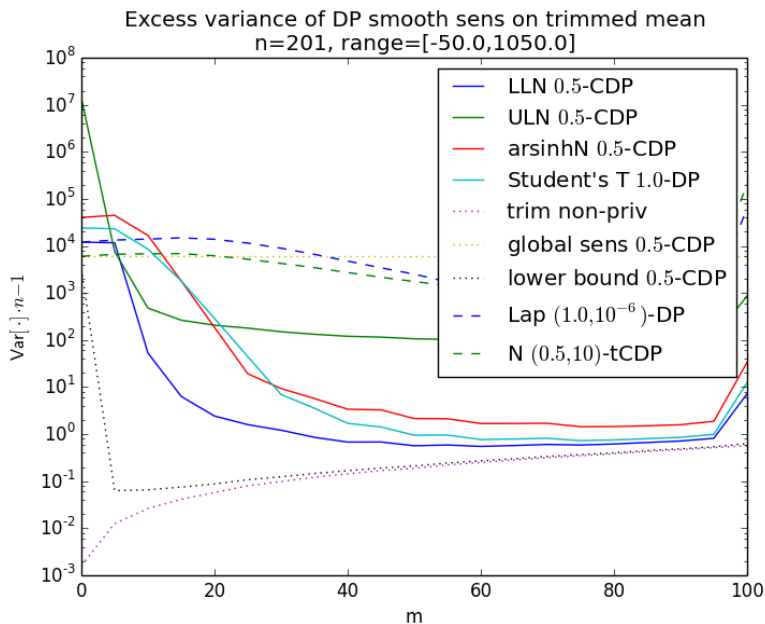
Lemma 20. *Let $Z \leftarrow \text{LLN}(\sigma)$ for $\sigma > 0$. Let $s \in \mathbb{R}$ and $\alpha \in (1, \infty)$. Then*

$$D_\alpha(Z \| Z + s) \leq \min \left\{ \frac{1}{2} e^{3\sigma^2} s^2 \alpha, e^{\frac{3}{2}\sigma^2} s \right\}.$$

- **Proof:**

Improvement from running smoothed sensitivity is substantial!

$$\text{trim}_m(x) = \frac{x_{(m+1)} + x_{(m+2)} + \dots + x_{(n-m)}}{n - 2m},$$



Bun and Steinke (2019): "Average case averages": <https://arxiv.org/pdf/1906.02830.pdf>

Next lecture

- Propose-Test-Release
- Stability-based query release
- Application to PATE