

Appendix: Examples of UNISEX Sessions

This appendix contains three pedagogical examples session to demonstrate the use of the UNISEX system. These transcripts are unedited except for setting user input in boldface.

The first example demonstrates the use of the system for testing a program using numeric values, a mix of numeric and symbolic values, and only symbolic values.

```
% unisex power.p
compiling power.p
loading unisex.....

1  program power;
2  var
3   a,b,raise,temb: integer;
4  begin
5   temb := b;
6   raise := 1;
7   while temb>0 do
8     begin
9       raise := raise*a;
10      temb := temb-1
11    end
12 end.
```

Execution mode ?

Enter (t)est or (v)erify: t

initialize debug functions

go

initialize variables

\$ **a=3, b=5**

pascal program execution completed

vars

Local variables:

a = 3

b = 5

raise = 243

temb = 0

pc: true

{ **Note that if tracing is not enabled very little is learned about the
internals of the program. 3 raised to the 5th power is 243 as expected.
Let's enable the branch option and try raising a negative number. }**

restart

```
1  program power;
```

Execution mode ?

Enter (t)est or (v)erify: t

initialize debug functions

branch on

```

# go
initialize variables
$ a=-13, b=3

predicate on line 7 evaluates to 3 > 0
true path taken

predicate on line 7 evaluates to 2 > 0
true path taken

predicate on line 7 evaluates to 1 > 0
true path taken

predicate on line 7 evaluates to 0 > 0
false path taken

pascal program execution completed
# vars
Local variables:
a = -13
b = 3
raise = -2197
temb = 0

pc: true
# { Again the correct result was obtained, but very little additional
  information was provided. }

# restart

1  program power;

Execution mode ?
Enter (t)est or (v)erify: t
initialize debug functions
# { To avoid having to restart each time we'll put a break at statement 4.
  We'll also enable full tracing by using the verbose option. Let's try
  a symbolic and a numeric value. }

# setbr 4
# verb on
# go
1  program power;
2  var
3  a,b,raise,temb: integer;
initialize variables
$ a=VALa, b=3
b = 3
a = VALa
Local variables:
a = VALa
b = 3
raise = undef3
temb = undef2

```

pc: true

breakpoint at line 4

```
4 begin
# { We'll save this for a starting point. }
```

save starting point

newly saved state -- 1 -- starting point

```
# go
5   temb := b;
temb = 3
6   raise := 1;
raise = 1
7   while temb>0 do
```

predicate on line 7 evaluates to $3 > 0$

true path taken

```
8   begin
9   raise := raise*a;
raise = VALa
10  temb := temb-1
temb = 2
11  end
7   while temb>0 do
```

predicate on line 7 evaluates to $2 > 0$

true path taken

```
8   begin
9   raise := raise*a;
raise = VALa*VALa
10  temb := temb-1
temb = 1
11  end
7   while temb>0 do
```

predicate on line 7 evaluates to $1 > 0$

true path taken

```
8   begin
9   raise := raise*a;
raise = VALa*VALa*VALa
10  temb := temb-1
temb = 0
11  end
7   while temb>0 do
```

predicate on line 7 evaluates to $0 > 0$

false path taken

```
12 end.
```

pascal program execution completed

```
# { The final value of raise contains the desired result. Now let's try
all symbolic values. First we restore the start state, then we set
variable values using the change command. }
```

```
# restore 1
state 1 restored -- starting point
next line to be executed -- 4
```

```
breakpoint at line 4
4 begin
# change a=symbolA
a = symbolA
# change b=symbolB
b = symbolB
# vars
Local variables:
a = symbolA
b = symbolB
raise = undef3
temb = undef2
```

```
pc: true
# go
5   temb := b;
temb = symbolB
6   raise := 1;
raise = 1
7   while temb>0 do
```

predicate on line 7 evaluates to symbolB > 0

```
-----
Does Path Condition --
true
```

```
-----
imply
symbolB > 0
-----
```

Enter (t)true, (f)false or (n)either

```
# { We have no way of telling if symbolB > 0, so we would like to try
both branches. Thus we save the current state. }
```

```
# n
Neither is implied.
ASSUME (t)true or (f)false
# save for false 1st time at loop
newly saved state -- 2 -- for false 1st time at loop
# t
New pc: symbolB > 0
```

```
-----
true path taken
8   begin
9   raise := raise*a;
raise = symbolA
10  temb := temb-1
temb = symbolB - 1
11  end
7   while temb>0 do
```

predicate on line 7 evaluates to $\text{symbolB} - 1 > 0$

Does Path Condition --

$\text{symbolB} > 0$

 imply

$\text{symbolB} > 1$

Enter (t)true, (f)false or (n)either

{ **We're at the loop again and again do not know the answer.** }

save for false 2nd time at loop

newly saved state -- 3 -- for false 2nd time at loop

n

Neither is implied.

ASSUME (t)true or (f)false

t

New pc: $\text{symbolB} > 1$

true path taken

 8 begin

 9 raise := raise*a;

raise = $\text{symbolA} * \text{symbolA}$

 10 temb := temb-1

temb = $\text{symbolB} - 2$

 11 end

 7 while temb>0 do

predicate on line 7 evaluates to $\text{symbolB} - 2 > 0$

Does Path Condition --

$\text{symbolB} > 1$

 imply

$\text{symbolB} > 2$

Enter (t)true, (f)false or (n)either

{ **Looks like more of the same. Let's save the current state and pursue some other paths to see what happens.** }

save third time at loop

newly saved state -- 4 -- third time at loop

sates

input not understood -- sates

states

currently stored states:

state	line	in	comment
-------	------	----	---------

4	7	power	third time at loop
---	---	-------	--------------------

3	7	power	for false 2nd time at loop
---	---	-------	----------------------------

2	7	power	for false 1st time at loop
---	---	-------	----------------------------

1	4	power	starting point
---	---	-------	----------------

restore 2

state 2 restored -- for false 1st time at loop

next line to be executed -- 7

```

7   while temb>0 do

predicate on line 7 evaluates to symbolB > 0
-----
Does Path Condition --
true
-----
    imply
symbolB > 0
-----
Enter (t)rue, (f)alse or (n)either
# n
Neither is implied.
ASSUME (t)rue or (f)alse
# f
New pc: symbolB <= 0
-----
false path taken
12 end.

pascal program execution completed
# vars
Local variables:
a = symbolA
b = symbolB
raise = 1
temb = symbolB

pc: symbolB <= 0
# { By looking at the pc we see that this path is followed if the
input value for b is less than or equal to zero. Thus, the result
is 1 whenever the input for b is non-positive. This may or may not
be what is desired. }

# restore 3
state 3 restored -- for false 2nd time at loop
next line to be executed -- 7
    7   while temb>0 do

predicate on line 7 evaluates to symbolB - 1 > 0
-----
Does Path Condition --
symbolB > 0
-----
    imply
symbolB > 1
-----
Enter (t)rue, (f)alse or (n)either
# n
Neither is implied.
ASSUME (t)rue or (f)alse
# f
New pc: symbolB = 1
-----

```

false path taken
12 end.

pascal program execution completed

vars

Local variables:

a = symbolA

b = symbolB

raise = symbolA

temb = symbolB - 1

pc: symbolB = 1

{ This looks good. Let's try one more iteration of the loop. }

restore 4

state 4 restored -- third time at loop

next line to be executed -- 7

7 while temb>0 do

predicate on line 7 evaluates to symbolB - 2 > 0

Does Path Condition --

symbolB > 1

imply

symbolB > 2

Enter (t) rue, (f)alse or (n)either

n

Neither is implied.

ASSUME (t) rue or (f)alse

f

New pc: symbolB = 2

false path taken

12 end.

pascal program execution completed

vars

Local variables:

a = symbolA

b = symbolB

raise = symbolA*symbolA

temb = symbolB - 2

pc: symbolB = 2

{ It appears that the program is working given that we want numbers raised to a negative power to yield the result 1. }

go

return to unix?

Enter (y)es or (n)o: y

%

The second example demonstrates the use of the UNISEX system to determine the restrictions that must be placed on symbolic values to cause a particular path to be traversed.

```
% unisex mpy.p
compiling mpy.p
loading unisex.....

1  program multiply;
2  var
3    a,b,prod,tema,temb,s: integer;
4  begin
5    tema := a;
6    prod := 0;
7    while tema<>0 do
8      begin
9        if tema>0 then
10         s := 1
11       else
12         s := -1;
13       temb := b;
14       while temb<>0 do
15         if temb>0 then
16           begin
17             prod := prod+s;
18             temb := temb-1
19           end
20         else
21           begin
22             prod := prod-s;
23             temb := temb+1
24           end;
25         tema := tema-s
26       end
27 end.
```

Execution mode ?

Enter (t)est or (v)erify: t

initialize debug functions

verb on

go

```
1  program multiply;
2  var
3    a,b,prod,tema,temb,s: integer;
```

initialize variables

\$ **a=A,b=B**

b = B

a = A

Local variables:

a = A

b = B

prod = undef5

tema = undef4

temb = undef3

s = undef2

```
pc: true
  4  begin
    5  tema := a;
tema = A
    6  prod := 0;
prod = 0
    7  while tema<>0 do
```

predicate on line 7 evaluates to $A \neq 0$

Does Path Condition --
true

imply
 $A \neq 0$

Enter (t)true, (f)false or (n)either

**# { Respond true for one iteration of outer loop. Must always respond
with neither first so that assumption is conjoined to pc. }**

n

Neither is implied.

ASSUME (t)true or (f)false

t

New pc: $A \neq 0$

true path taken

```
    8  begin
    9  if tema>0 then
```

predicate on line 9 evaluates to $A > 0$

Does Path Condition --
 $A \neq 0$

imply
 $A > 0$

Enter (t)true, (f)false or (n)either

n

Neither is implied.

ASSUME (t)true or (f)false

f

New pc: $A < 0$

false path taken

```
    11  else
    12  s := -1;
s = -1
    13  temb := b;
temb = B
    14  while temb<>0 do
```

predicate on line 14 evaluates to $B \leq 0$

Does Path Condition --
 $A < 0$

 imply
 $B \leq 0$

Enter (t)true, (f)false or (n)either
n
Neither is implied.
ASSUME (t)true or (f)false
t
New pc: $A < 0$ and $B \leq 0$

true path taken
 15 if $temb > 0$ then

predicate on line 15 evaluates to $B > 0$

Does Path Condition --
 $A < 0$ and $B \leq 0$

 imply
 $B > 0$

Enter (t)true, (f)false or (n)either
n
Neither is implied.
ASSUME (t)true or (f)false
t
New pc: $B > 0$ and $A < 0$

true path taken
 16 begin
 17 $prod := prod + s;$
 $prod = -1$
 18 $temb := temb - 1$
 $temb = B - 1$
 19 end
 14 while $temb \leq 0$ do

predicate on line 14 evaluates to $B - 1 \leq 0$

Does Path Condition --
 $B > 0$ and $A < 0$

 imply
 $B > 1$

Enter (t)true, (f)false or (n)either
n
Neither is implied.
ASSUME (t)true or (f)false

t

New pc: $B > 1$ and $A < 0$

true path taken

15 if $temb > 0$ then

predicate on line 15 evaluates to $B - 1 > 0$

true path taken

16 begin

17 $prod := prod + s;$

$prod = -2$

18 $temb := temb - 1$

$temb = B - 2$

19 end

14 while $temb \leq 0$ do

predicate on line 14 evaluates to $B - 2 \leq 0$

Does Path Condition --

$B > 1$ and $A < 0$

imply

$B > 2$

Enter (t)true, (f)false or (n)either

{ **Respond false to exit after second iteration.** }

n

Neither is implied.

ASSUME (t)true or (f)false

f

New pc: $B = 2$ and $A < 0$

false path taken

25 $tema := tema - s$

$tema = A + 1$

26 end

7 while $tema \leq 0$ do

predicate on line 7 evaluates to $A + 1 \leq 0$

Does Path Condition --

$B = 2$ and $A < 0$

imply

$A < -1$

Enter (t)true, (f)false or (n)either

{ **Respond false to exit outer loop after one iteration.** }

n

Neither is implied.

ASSUME (t)true or (f)false

f

New pc: A = -1 and B = 2

false path taken
27 end.

pascal program execution completed

{ **Current value of pc gives restrictions on input values needed to force this path to be followed. In this case the values are restricted to a single possible value for A, -1, and for B, 2. The current path can be seen b using the paths command. }**

paths

current path: 1,2,3,4,5,6,7,8,9,11,12,13,14,15,16,17,18,19,14,15,16,17,18,19,14,25,26,7,27

restart

1 program multiply;

Execution mode ?

Enter (t)est or (v)erify: t

initialize debug functions

{ **We will now try input values -1 and 2. }**

go

initialize variables

\$ **a=-1, b=2**

pascal program execution completed

vars

Local variables:

a = -1

b = 2

prod = -2

tema = 0

temb = 0

s = -1

pc: true

paths

current path: 1,2,3,4,5,6,7,8,9,11,12,13,14,15,16,17,18,19,14,15,16,17,18,19,14,25,26,7,27

q

return to unix?

Enter (y)es or (n)o: y

%

The final example demonstrates how to generate the necessary verification conditions to show that a program is consistent with its specifications.

```
% unisex divide.p
compiling divide.p
loading unisex.....
```

```
1  program divide(input,output);
2  var
3    x,y,quot,rem: integer;
4  { : entry ( (x>=0),(y>=0) ) : }
5  { : exit ( (x'=quot*y'+rem),(rem>=0),(rem<y') ) : }
6  begin
7    quot := 0;
8    rem := x;
9    { : assert ((x=quot*y+rem),(rem>=0),(x=x'),(y=y')) : }
10 while rem>=y do
11   begin
12     quot := quot+1;
13     rem := rem-y
14   end
15 end.
```

Execution mode ?

Enter (t)est or (v)erify:

v

initialize debug functions

{ **This program illustrates the use of an assert-while loop, as well as entry and exit specifications.** }

verb on

go

```
1  program divide(input,output);
2  var
3    x,y,quot,rem: integer;
initializing variables:
4  { : entry ( (x>=0),(y>=0) ) : }
5  { : exit ( (x'=quot*y'+rem),(rem>=0),(rem<y') ) : }
New pc: $x >= 0 and $y >= 0
6  begin
7    quot := 0;
quot = 0
8    rem := x;
rem = $x
9    { : assert ((x=quot*y+rem),(rem>=0),(x=x'),(y=y')) : }
```

predicate on line 9 evaluates to:

\$y = \$y and \$x = \$x and \$x >= 0 and \$x = 0*\$y + \$x

** Prove Condition **

Does Path Condition --

\$x >= 0 and \$y >= 0

```

-----
    imply
true
-----

    Path verified
-----
initializing variables:
New pc: @quot*@y + @rem = @x and $x = @x and $y = @y and @rem >= 0
    10 while rem>=y do

predicate on line 10 evaluates to @rem >= @y
-----
Does Path Condition --
@quot*@y + @rem = @x and $x = @x and $y = @y and @rem >= 0
-----
    imply
@rem >= @y
-----
Enter (t)true, (f)false or (n)either
# { After the Verification Condition (VC) was taken care of by the simplifier,
    UNISEX reinitialized all of the variables. The current path is not related
    to the one just completed. Here we have an unresolvable decision. Unisex
    prompts the user (as theorem prover) to make a decision. In automatic mode
    a decision of “neither” would result in the false path being stacked
    and the true path being continued. }

# n
Neither is implied.
ASSUME (t)true or (f)false
# { Since there are two paths starting at this point we save the current state. }

# save entry to exit (for false branch)
newly saved state -- 1 -- entry to exit (for false branch)
# t
New pc: @quot*@y + @rem = @x
        and $x = @x
        and $y = @y
        and @rem >= 0
        and @rem >= @y
-----
true path taken
    11 begin
    12 quot := quot+1;
quot = @quot + 1
    13 rem := rem-y
rem = @rem - @y
    14 end
    9 { : assert ((x=quot*y+rem),(rem>=0),(x=x'),(y=y')) : }

predicate on line 9 evaluates to:
    @y = $y and @x = $x and @rem - @y >= 0 and @x = (@quot + 1)*@y + @rem - @y
-----
** Prove Condition **
-----

```

Does Path Condition --

@quot*@y + @rem = @x and \$x = @x and \$y = @y and @rem >= 0 and @rem >= @y

imply

true

Path verified

initializing variables:

New pc: @quot*@y + @rem = @x and \$x = @x and \$y = @y and @rem >= 0

10 while rem>=y do

predicate on line 10 evaluates to @rem >= @y

Does Path Condition --

@quot*@y + @rem = @x and \$x = @x and \$y = @y and @rem >= 0

imply

@rem >= @y

Enter (t)rue, (f)alse or (n)either

paths

current path: 9,10

Paths verified:

9,10,11,12,13,14,9

3,4,5,6,7,8,9

No unverifiable paths.

{ **At this point 2 of the 3 paths in the program have been verified. The decision we are being asked to make now is exactly the same as the last one.** }

states

currently stored states:

state	line	in	comment
-------	------	----	---------

1	10	divide	entry to exit (for false branch)
---	----	--------	----------------------------------

restore 1

state 1 restored -- entry to exit (for false branch)

next line to be executed -- 10

10 while rem>=y do

predicate on line 10 evaluates to @rem >= @y

Does Path Condition --

@quot*@y + @rem = @x and \$x = @x and \$y = @y and @rem >= 0

imply

@rem >= @y

Enter (t)rue, (f)alse or (n)either

n

Neither is implied.

ASSUME (t)rue or (f)alse


```
# { Note that UNISEX does not remember that we told it this decision was
  unresolvable before we saved the state. We have to tell it again. }
```

```
# f
```

```
New pc:  @quot*@y + @rem = @x
         and $x = @x
         and $y = @y
         and @rem >= 0
         and @rem < @y
```

```
-----
false path taken
15 end.
```

predicate on line 15 evaluates to:

```
@rem < $y and @rem >= 0 and $x = @quot*$y + @rem
```

```
-----
** Prove Condition **
```

```
-----
Does Path Condition --
```

```
@quot*@y + @rem = @x and $x = @x and $y = @y and @rem >= 0 and @rem < @y
```

```
-----
imply
```

```
$x - $y*@quot = @rem and $y > @rem
```

```
-----
Enter (t)true, (f)false or (n)either
```

```
# { This VC shows one limitation of the simplifier; it does not substitute
  for equal values. We can replace $x with @x and $y with @y and see that
  the VC is indeed true. }
```

```
# t
```

```
Path verified
```

```
-----
pascal program execution completed
```

```
# paths
```

```
current path: 9,10,15
```

```
Paths verified:
```

```
9,10,15
```

```
9,10,11,12,13,14,9
```

```
3,4,5,6,7,8,9
```

```
No unverifiable paths.
```

```
# { The verification of program divide is complete (assuming that the
  user-theorem prover is sound!). Here are some other commands: }
```

```
# help
```

```
111.
```

```
go quit restart
```

```
list list n list n1,n2
```

```
vars var varname change var=value
```

```
states pc types
```

```

save name  restore number  rmstate number
curbr setbr n1,...  rmbr n1,...
brk on|off  verb on|off  step on|off
branch on|off  paths  addpred (expr)
simplify on|off  axioms  verify routine
autoverify routine  autoverify program

# list 9
9  { : assert ((x=quot*y+rem),(rem>=0),(x=x'),(y=y')) : }

# list 2,8
2  var
3  x,y,quot,rem: integer;
4  { : entry ( (x>=0),(y>=0) ) : }
5  { : exit ( (x'=quot*y'+rem),(rem>=0),(rem<y') ) : }
6  begin
7  quot := 0;
8  rem := x;

# curbr
No breakpoints set.

# setbr 12
# curbr
Current breakpoints:
12

# q
return to unix?
Enter (y)es or (n)o: y

%
```