
CS-235
Computational Geometry

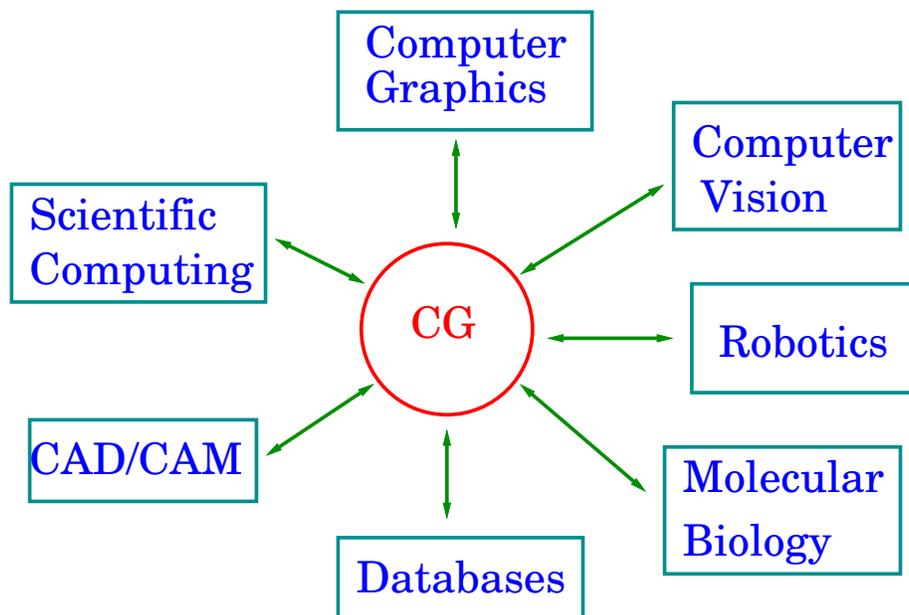
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Computational Geometry

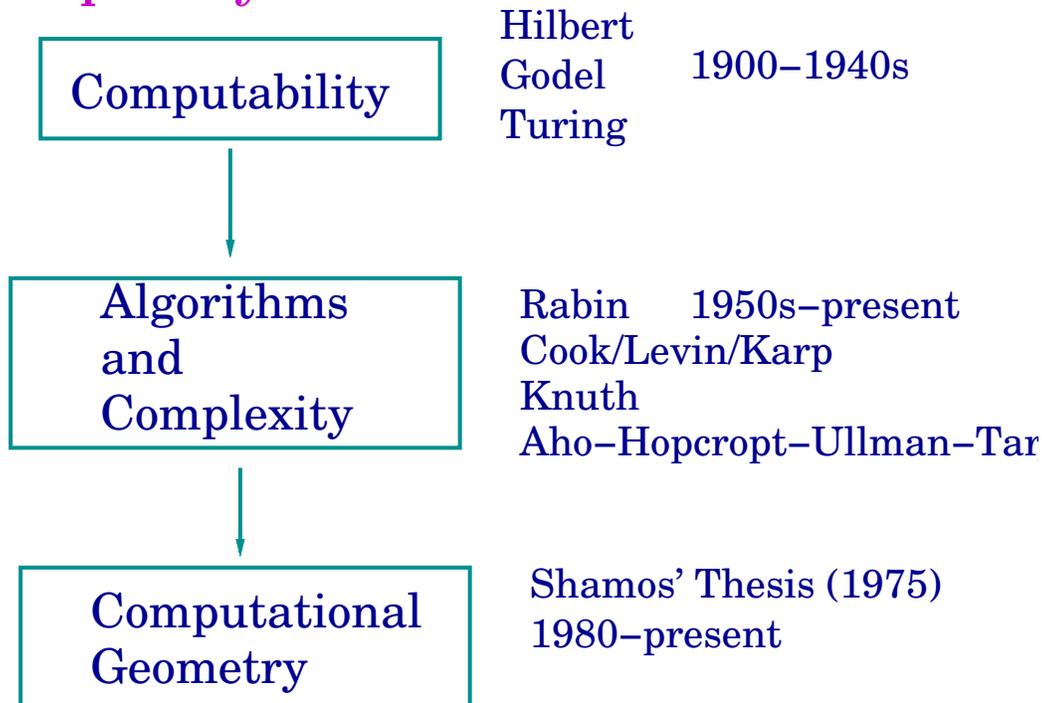
- Study of algorithms for geometric problems.
- Deals with **discrete** shapes: points, lines, polyhedra, polygonal meshes.
- **Abstraction** of problems in different applied areas.



- Occlusion, visibility, augmented reality, collision detection, motion or assembly planning, drug design, databases, GIS, layout, fluid dynamics, etc.

CG and Computer Science

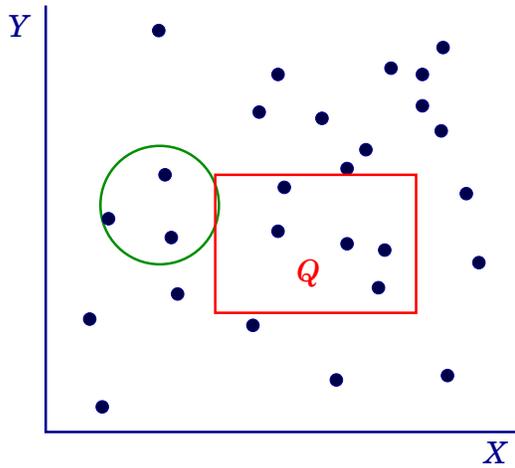
- **CG is a sub-discipline of algorithms and complexity.**



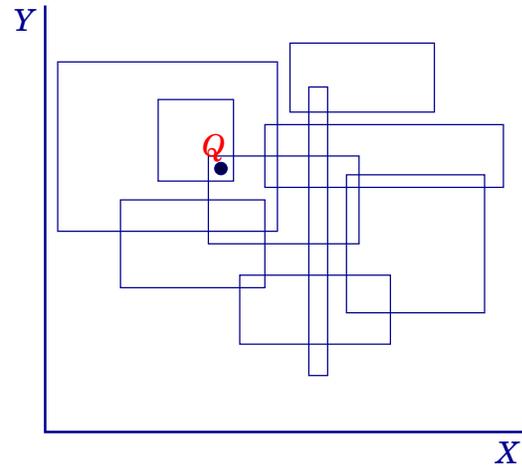
- Develops **fundamental techniques and tools** for geometric problems.
- **Motivated** by applications in other CS fields.
- **Significant** interaction with discrete mathematics.

Some Examples

- Range Searching Data Structures.

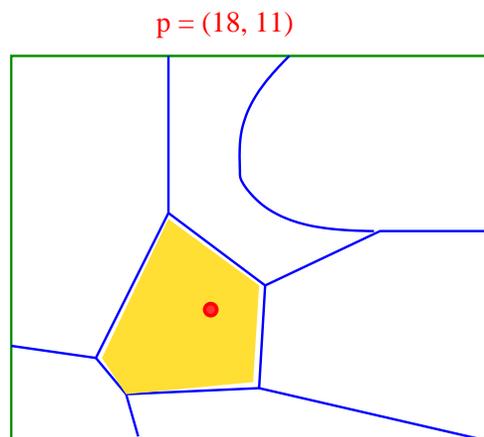


Point objects



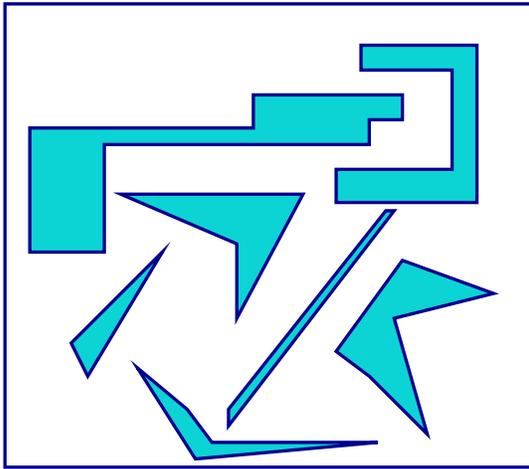
Rectangular objects

- Location Queries.

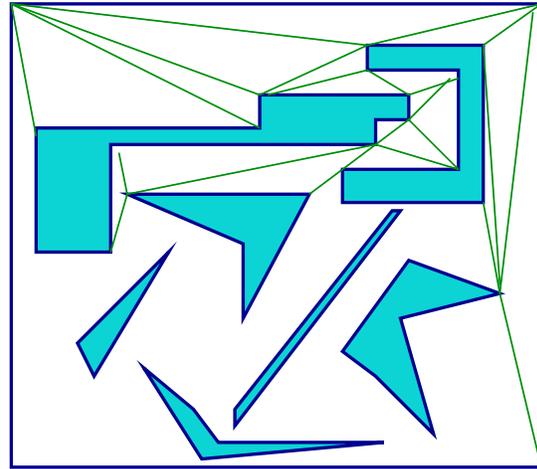


Some Examples

- **Decomposition.**



Geometric Scene

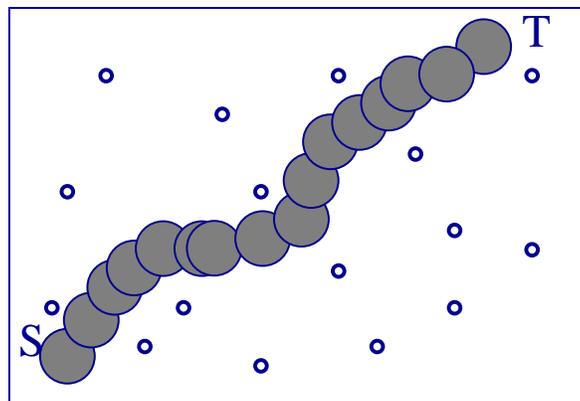
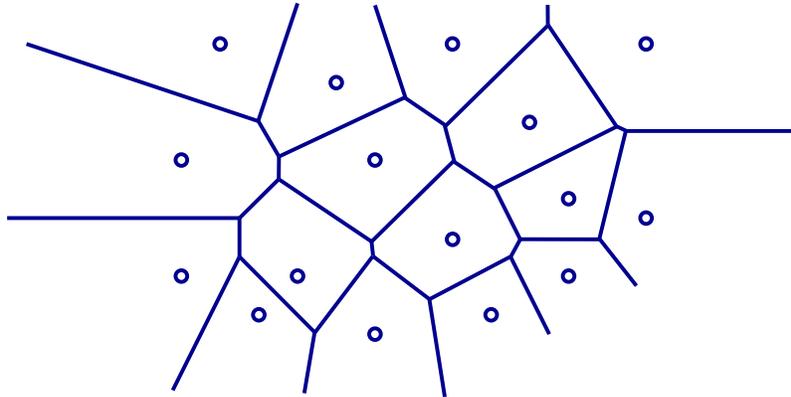


Partial Triangulation

- Is this always possible?
- In three dimension?
- Other examples: Shortest paths, geometric structures, visibility, pattern matching.

Some Examples

- Spatial Data Structures.



- Voronoi diagram, Delaunay triangulation.
- Robot motion planning.

Taste of Comb. Geometry

- **Helly's Theorem:** Let C_1, \dots, C_n be a family of convex sets in the plane. If every triple intersects, then $\bigcap C_i$ is non-empty.
- **Center Points:** Given points p_1, p_2, \dots, p_n in the plane, a point x is called **center point** if any line through x contains **at least $n/3$** points on each side.
- **Ham Sandwich Theorem:** Take n red points and n blue points in the plane. There is a line **simultaneously bisecting** both red and blue points.
- **Crossing Number Theorem:** If G is a graph with n nodes and m edges, then **every drawing** of G in the plane contains at least $c \left(\frac{m^3}{n^2} \right) - n$ crossings.

Taste of Comb. Geometry

- Among any 5 points in the plane in general position, we can find 4 forming a convex polygon.
- **Erdős-Szekeres Theorem:** For every positive integer k , there exists a number F_k , such that every set of F_k points in the plane contains k that form a convex k -gon.
- **Empty k -gon:** How large must the set be to guarantee that we can find k points forming a convex polygon, which does not include any other point inside?
- **Empty k -gon: Known values:**
 - $G_3 = 3.$
 - $G_4 = 5.$
 - $G_5 = 10.$
 - $G_6 = ???.$
 - $G_7 = \infty!$

Spirit of CG

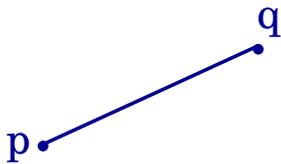
1. **CG** is a product of marriage between classical geometry and computer science.
2. **Emphasis on design of efficient algorithms and data structures.**
3. **In classical approach**, reducing the complexity to a **finite** number of choices was enough.
4. Alas! 10^{100} is mathematically finite but **computationally infinite**.
5. Point of Reference: **1 Year** $\approx 3.15 \times 10^7$ seconds.
6. **1 Century** $\approx \pi \times 10^9$ seconds.
7. **A 1-Giga flop computer does only** 10^{20} ops in a century!

Model of Computation

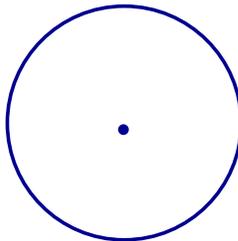
1. Assume an abstract programming language model.
2. Traditionally, real-number extension of Random Access Machine (RAM).
3. Each memory cell can hold one int or real geometric coordinate. We will discuss later numerical precision issues.
4. Standard repertoire of operators:
 - Arithmetic: $+$, $-$, \div , $*$, $\sqrt{\cdot}$.
 - Trigonometry: \sin , \cos , \tan , \exp , \log .
 - Comparators: \leq , \geq , $=$.
 - Array indexing, pointers.

Elementary Objects

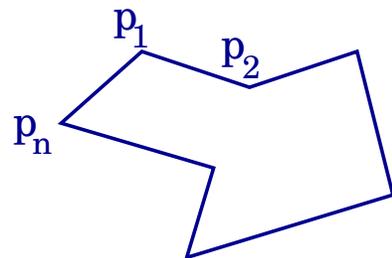
1. **Point** $p = (x, y)$, where x, y reals.
2. **Line** $\ell := ax + by = 1$.
3. **Line segment** $s = [p, q]$.
4. **Circle** $C = (p, r)$. (Center, radius)
5. **Polygon** $\langle p_1, p_2, \dots, p_n \rangle$.



Line Segment



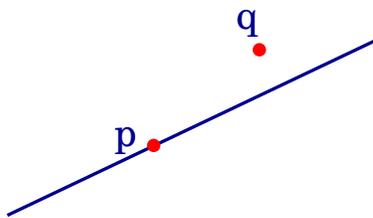
Circle



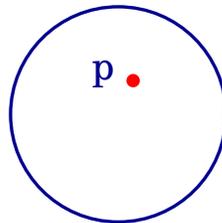
Polygon

Elementary Operations

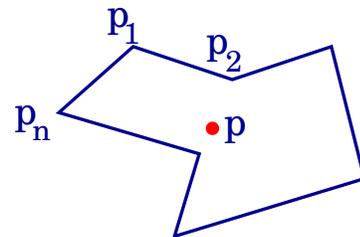
1. Algorithm receives **algebraic input**. It must perform **computation** to **see** the **underlying geometric relationships**.
2. **Is point p on line ℓ ?**
3. **Is point p inside or outside circle C ?**
4. **Do segments s_1 and s_2 intersect?**
5. **Is point p inside or outside polygon P ?**



Point on Line



Point in Circle



Point in Polygon

Overview of the Course

1. Convex Hulls.
2. Intersection Detection and Reporting.
3. Triangulation.
4. Range Searching.
5. Point Location.
6. Delaunay Triangulation.
7. Voronoi Diagrams.
8. Arrangements.
9. Binary Space Partitions.
10. Epsilon Net and VC Dimension.
11. Volume and paradoxes in higher dimensions.
12. Misc.