

Lecture 2: *August 8th, 2022*Lecturer: *Tali Kaufman**Intro to High Dimensional Expanders and Local to Global theorems*

These notes were prepared by Tushant Mittal based on the pair of lectures by Tali Kaufman. This is part of the 2022 Summer School on *New tools for optimal mixing of Markov chains: Spectral independence and entropy decay*, which was held at the University of California, Santa Barbara (UCSB) from August 8, 2022 to August 12, 2022. More information on the summer school is available at: <https://sites.cs.ucsb.edu/~vigoda/School/>

## 2.1 Introduction

In these two lectures, we will introduce *simplicial complexes* and prove two (local-to-global) results about them. Namely,

1. Trickleing Down Theorem - Bound the spectral gap of the local walks on any link using the local walks on  $(d - 2)$ -links (here  $d$  is the dimension of the complex).
2. Random Walk Theorem - Using the spectral gap of local walks for all links to bound the spectral gap of global walks, such as the Glauber dynamics.

## 2.2 Basic Definitions

A *simplicial complex* is a downward closed set system. It is called *pure* if every set is contained in a subset of maximal size. The *dimension* of the complex is  $d$  if the largest set in the system has size  $d + 1$ . We denote  $X(k)$  to be the sets of size exactly  $k + 1$  and define  $X(-1) = \{\phi\}$ . The *link* of a complex is related to the notion of conditioning. Formally,

**Definition 2.1** (Link). *Let  $\tau \in X(k)$ , then the link at  $\tau$  is a  $(d - k - 1)$ -dimensional simplicial complex,  $X_\tau = \{\sigma \setminus \tau \mid \tau \subset \sigma\}$ .*

**Weights:** Let  $w : X(d) \rightarrow \mathbb{R}_{\geq 0}$  be a probability measure on the top faces, i.e.,  $\sum_{\sigma \in X(d)} w(\sigma) = 1$ . We can define a measure on the lower levels  $X(k)$  by trickling the weights down,

$$\forall \tau \in X(k), \quad w(\tau) = \frac{1}{k+2} \sum_{\substack{\sigma \in X(k+1), \\ \tau \subset \sigma}} w(\sigma).$$

One can also define the measure on the links by taking marginals as follows,

$$\forall \sigma \in X_\tau, \quad w_\tau(\sigma) = \frac{w(\tau \cup \sigma)}{\binom{|\tau|+|\sigma|}{|\tau|} w(\tau)}.$$

**Cochains:** The space of  $k$ -cochains denoted as  $C^k(X, \mathbb{R}) := \{f : X(k) \rightarrow \mathbb{R}\}$  is a vector space with an inner product defined as  $\langle f, g \rangle_w = \sum_{\sigma \in X(k)} w(\sigma) f(\sigma) g(\sigma)$ . We will denote it simply as  $C^k$  when there is no ambiguity in the underlying complex.

**Definition 2.2** (Up-Down Operators). For each  $k < d$ , we define the following pair of adjoint operators between  $C^k$  and  $C^{k+1}$ .

(Up-Operator)  $P_k^\uparrow : C^k \rightarrow C^{k+1}$ . Let  $f$  be a function that is defined on  $(k+1)$ -face, i.e.,  $f \in C^k$ . The up operator lifts it to a function on  $(k+2)$ -faces by averaging over all subsets of size  $k+1$ . More formally,

$$P_k^\uparrow f(\sigma) = \mathbb{E}_{\substack{\tau \sim X(k) \\ \tau \subset \sigma}} [f(\tau)] = \frac{1}{k+2} \sum_{\substack{\tau \sim X(k) \\ \tau \subset \sigma}} f(\tau)$$

(Down-Operator)  $P_k^\downarrow : C^{k+1} \rightarrow C^k$ . Analogously, one can transform a function on  $k+2$  faces to one on  $k+1$  faces by averaging over all  $(k+2)$ -sized faces containing a given  $(k+1)$ -face.

$$P_k^\downarrow f(\tau) = \mathbb{E}_{\substack{\sigma \sim X(k+1) \\ \tau \subset \sigma}} [f(\sigma)] = \sum_{\substack{\sigma \in X(k+1) \\ \tau \subset \sigma}} w_\sigma(\sigma) f(\sigma)$$

By composing these two operators appropriately, we can define two kinds of walks at level  $k$ .

- (Up-Down Chain) This walk corresponds to the following random process. Start from  $\tau \in X(k)$ . In the up-step, we sample a random  $j$  such that  $\tau \cup \{j\} \in X(k+1)$ ; the probability of picking  $j$  is proportional to  $w(\tau \cup \{j\})$ . Then in the down step, we drop a uniform element of  $\tau \cup \{j\}$ .

$$P_k^\wedge : C^k \rightarrow C^k, \quad P_k^\wedge = P_{k+1}^\downarrow P_{k+1}^\uparrow.$$

This walk has a lazy component as there is a  $\frac{1}{k+2}$  chance of returning to  $\tau$ . To avoid this, we define the non-lazy version  $\tilde{P}_k^\wedge$  by subtracting the identity component,

$$\tilde{P}_k^\wedge : C^k \rightarrow C^k, \quad P_k^\wedge = \frac{k+1}{k+2} \tilde{P}_k^\wedge + \frac{1}{k+2} I.$$

- (Down-Up Chain) Similarly, we can first remove a uniformly random element  $i$  from  $\tau$  first and then add an element  $j$  with probability proportional to the weight of the resulting set  $w(\tau \cup \{j\} \setminus \{i\})$ ,

$$P_k^\vee : C^{k+1} \rightarrow C^{k+1}, \quad P_k^\vee = P_{k+1}^\uparrow P_{k+1}^\downarrow.$$

*Remark 2.3* (Glauber dynamics). For a spin system, the down-up chain  $P_{n-1}^\vee$  for the appropriate simplicial complex is equivalent to the Glauber dynamics. A spin system is defined on a graph  $G = (V, E)$  and the Gibbs distribution  $\mu$  has support  $\Omega \subset \{1, \dots, q\}^V$  for integer  $q \geq 2$ . The elements of the corresponding simplicial complex are (vertex, spin) pairs  $(v, \sigma(v))$  where  $v \in V$  and  $\sigma(v) \in \{1, \dots, q\}$ .

Let us illustrate the above for the special case of independent sets (this is the hard-core model with  $\lambda = 1$ ); here the Gibbs distribution is uniformly distributed over all independent sets (of any size) of  $G$ . The corresponding simplicial complex has dimension  $n = |V|$ ; for every independent set  $I$  we have  $\{(v, 1) \mid v \in I\} \cup \{(v, 0) \mid v \notin I\}$  in  $X(n-1)$ . The Glauber dynamics which updates the spin at a randomly chosen vertex is equivalent to  $P_{n-1}^\vee$ .

We are now ready to define the notion of spectral expansion for a simplicial complex. Let  $\tau \in X(k)$  for  $k \leq d-2$ . The 1-skeleton of  $X_\tau$  is the weighted graph  $G_\tau = (X_\tau(0), X_\tau(1))$ . The weight of vertices (or edges) in  $G_\tau$  are obtained by taking the weights in  $X$  and dividing by  $w(\tau)$ . Note that the sum of the weights of all the vertices (or edges) in  $G_\tau$  is 1. The eigenvalues of  $G_\tau$  are solutions of  $Ax = \lambda Dx$  where  $A$  is the weighted adjacency matrix and  $D$  is a diagonal matrix with vertex weights; the weight of a vertex is the sum of the weights of the adjacent edges, this follows from the above definition of vertex/edge weights.

**Definition 2.4** ( $\lambda$ -local spectral expanders). *A pure  $d$ -dimensional simplicial complex  $X$  is a  $\lambda$ -local spectral expander if for every  $\tau \in X(k)$  such that  $k \leq d-2$ , the 1-skeleton of  $X_\tau$ , which is the weighted graph  $G_\tau = (X_\tau(0), X_\tau(1))$ , satisfies  $\lambda_2(G_\tau) \leq \lambda$ .*

**Local walk:** We refer to the random walk  $\tilde{P}_0^\wedge$  on the 1-skeleton of  $X_\tau$  as the *local walk*. Note that the spectrum of the local walk on  $X_\tau$  is the same as the spectrum of  $G_\tau$ .

## 2.3 The Trickle-down theorem

**Theorem 2.5** (Oppenheim [Opp18]). *If  $X$  is a pure simplicial complex such that*

- (i) *Its 1-skeleton is connected,*
- (ii)  $\forall v \in X(0), \lambda_2(X_v(0), X_v(1)) \leq \lambda,$

*then,  $X$  is a  $\frac{\lambda}{1-\lambda}$ -local spectral expander.*

**Corollary 2.6** (Trickle with loss). *If  $X$  is  $d$ -dimensional and all  $(d-2)$ -links are  $\lambda$ -expanders, then  $X$  is a  $\frac{\lambda}{1-(d-2)\lambda}$ -local spectral expander.*

**Corollary 2.7** (Trickle without loss). *If  $X$  is  $d$ -dimensional and all  $(d-2)$ -links are 0-local expanders, then  $X$  is a 0-local spectral expander.*

## 2.4 Proof of Trickle-Down

The goal is to show that if  $\lambda_2(X_v) \leq \lambda$  for every  $v \in X(0)$ , then, we have  $\lambda_2(X) \leq \frac{\lambda}{1-\lambda}$ . To do so, we will need a way to study functions locally and the key notion we will use here will be that of restriction.

**Restriction:** Let  $\tau \in X(i)$  and let  $f \in C^k$ . We define the restriction of  $f$  to  $X_\tau$ ,  $f^\tau \in C^k(X_\tau, \mathbb{R})$ , to be  $f^\tau(\sigma) = f(\sigma)$ . Note that here,  $\sigma \in X_\tau(k)$  which means that  $\sigma \cup \tau \in X(i+k+1)$  and by the downward closed property,  $\sigma \in X(k)$  and therefore  $f^\tau$  is well-defined.

We wish to bound the second largest eigenvalue of the (weighted) adjacency operator on the (global) graph  $G = (X(0), X(1))$ . This operator, which we denote as  $\tilde{P}_0^\wedge$ , can be seen as the non-lazy part of the up-down operator from vertices to edges and back to vertices. More formally, the up-down operator is  $P_0^\wedge = P_1^\downarrow P_1^\uparrow : C^0 \rightarrow C^0$  which decomposes as  $P_0^\wedge = \frac{1}{2} \tilde{P}_0^\wedge + \frac{1}{2} I$ .

**Lemma 2.8** (Restriction Lemma). *For a  $d$ -dimensional simplicial complex  $X$ , cochains  $f, g \in C^k(X, \mathbb{R})$  and  $0 \leq i \leq d-k-2$ , we have,*

1.  $\langle f, g \rangle = \mathbb{E}_{\tau \in X(i)} [\langle f^\tau, g^\tau \rangle],$

$$2. \langle \tilde{\mathbb{P}}_0^\wedge f, g \rangle = \mathbb{E}_{v \in X(0)} [\langle \tilde{\mathbb{P}}_{0,v}^\wedge f^v, g^v \rangle].$$

*Proof.* Exercise. □

Now, we will prove the trickle-down theorem. Let  $f \in C^0$  be an eigenfunction perpendicular to the constant function, that is,  $(M_0')^+ f = \mu f$  and  $\|f\| = 1$ . We will bound  $\mu$  in terms of the local spectral expansion  $\lambda$ .

*Proof of Trickle-down.*

$$\begin{aligned} \mu &= \langle \tilde{\mathbb{P}}_0^\wedge f, f \rangle \\ &= \mathbb{E}_{v \in X(0)} \left[ \langle \tilde{\mathbb{P}}_{0,v}^\wedge f^v, f^v \rangle \right] && \text{Restriction Lemma} \\ &= \mathbb{E}_{v \in X(0)} \left[ \left[ \langle \tilde{\mathbb{P}}_{0,v}^\wedge f^{v^\perp}, f^{v^\perp} \rangle + \langle \tilde{\mathbb{P}}_{0,v}^\wedge f^{v^\parallel}, f^{v^\parallel} \rangle \right] \right] \\ &\leq \mathbb{E}_{v \in X(0)} \left[ \left[ \lambda \|f^{v^\perp}\|^2 + \|f^{v^\parallel}\|^2 \right] \right] && \text{Using local expansion} \\ &\leq \lambda \mathbb{E}_{v \in X(0)} \left[ \|f^v\|^2 \right] + (1 - \lambda) \mathbb{E}_{v \in X(0)} \left[ \|f^{v^\parallel}\|^2 \right]. \quad \|f^{v^\perp}\|^2 = \|f^v\|^2 - \|f^{v^\parallel}\|^2 \end{aligned}$$

We can bound the second term using the following observation,

$$\|f^{v^\parallel}\| = |\langle f^v, \mathbf{1}_v \rangle| = \left| \mathbb{E}_{(u,v) \in X(1)} [f^v(u)] \right| = |\tilde{\mathbb{P}}_0^\wedge f(v)|. \quad (2.1)$$

Here,  $\mathbf{1}_v$  denotes the constant all-ones function in  $C^0(X_v, \mathbb{R})$ .

$$\begin{aligned} \mu &\leq \lambda \mathbb{E}_{v \in X(0)} [\langle f^v, f^v \rangle] + (1 - \lambda) \mathbb{E}_{v \in X(0)} \left[ |\tilde{\mathbb{P}}_0^\wedge f(v)|^2 \right] && \text{Using Eq. (2.1)} \\ &\leq \lambda \langle f, f \rangle + (1 - \lambda) \left\| \tilde{\mathbb{P}}_0^\wedge f \right\|^2 && \text{Using Restriction lemma} \\ &= \lambda \|f\|^2 + (1 - \lambda) \mu^2 \|f\|^2. \end{aligned}$$

Solving this finishes the proof,

$$\begin{aligned} \mu &\leq \lambda + (1 - \lambda) \mu^2 \\ \mu - \mu^2 &\leq \lambda(1 - \mu^2) \\ \mu &\leq \lambda(1 + \mu) \\ \mu &\leq \frac{\lambda}{1 - \lambda}. \end{aligned} \quad \square$$

## 2.5 Random Walk Theorem

Today we will give a proof of the spectral gap of the random walks. Recall the following definitions from the first lecture. Recall the up-down operators from [Definition 2.2](#). We wish to prove rapid mixing of the walks at level  $k$ .

**Theorem 2.9** (Convergence of RW in local-spectral expander [KM17, DK17, KO20]). *If  $X$  is a pure  $d$ -dimensional simplicial complex which is a  $\gamma$ -local spectral expander, then for any  $1 \leq k \leq d$ ,*

$$\lambda_2(P_k^\vee) \leq 1 - \frac{1}{k+1} + O(\gamma k).$$

**Theorem 2.10** ([AL20]). *If  $X$  is a  $(\lambda_{-1}, \dots, \lambda_{d-2})$ -local spectral expander where  $\gamma_j = 1 - \max_{\tau \in X(j)} \lambda_2(X_\tau)$ , then for any  $1 \leq k \leq d$ ,*

$$\lambda_2(P_k^\vee) = \lambda_2(P_k^\wedge) \leq 1 - \frac{1}{k+1} \prod_{i=-1}^{k-2} \gamma_i.$$

**Localization** Given a global function, we wish to define a local function on the links. Let  $f \in C^k, \sigma \in X(i)$ . We define  $f_\sigma : X_\sigma(k-i) \rightarrow \mathbb{R}$  as  $f_\sigma(\tau) = f(\sigma \cup \tau)$ .

*Example 2.11.* Let  $f \in C^1$ , i.e., a function on edges. For a vertex  $v$ , we have  $f_v(u) = f((u \cup v))$  where  $u \in X_v(0)$ . Therefore,  $f_v \in C^0(X_v, \mathbb{R})$

This differs from the restriction we saw yesterday as the restriction of  $f$  gives a local function  $f^v \in C(X_\tau^k)$  whereas the localization is  $f_v \in C(X_\tau^{k-i})$ . While restriction worked well with trickle down, localization is better suited for *Garland's method*.

**Lemma 2.12** (Localization). *Let  $f, g \in C^k$ . We have that,*

1.  $\langle f, g \rangle = \mathbb{E}_{\sigma \in X(i)} [\langle f_\sigma, g_\sigma \rangle],$
2.  $\langle \tilde{P}_k^\wedge f, f \rangle = \mathbb{E}_{v \in X(0)} [\langle \tilde{P}_{k-1,v}^\wedge f_v, f_v \rangle].$

*Proof.* Proof not provided in the lecture. □

**Properness:** A cochain  $f \in C^k$  is an  $i$ -level co-chain if  $f \in \ker(P_{i-1}^\downarrow \cdots P_{k-1}^\downarrow)$ . We denote this subspace as  $C_i^k$ . We say that  $f \in C_i^k$  is a *proper  $i$ -level co-chain* if  $f \in \text{Im}(P_{k-1}^\uparrow \cdots P_i^\uparrow)$ . Essentially, every function at level  $i$  is also in  $i-1$ . Whereas a proper level  $i$  function is one not in level  $i+1$ .

**Lemma 2.13** (Orthogonal Decomposition). *Every  $f \in C^k$  can be represented as  $f = \sum_i f^i$  such that  $f^i$  is proper  $i$ -level cochain and  $\langle f^i, f^j \rangle = 0$  whenever  $i \neq j$ .*

*Proof.* Proof not given in the lecture. □

Before we prove the main theorem, we make the following observation about the localization of a cochain  $f$ .

**Claim 2.14** (Localization Properties). *Let  $f \in C^k$  with the decomposition  $f = \sum_i f^i, f^i \in C_i^k$  and  $v \in X(0)$ . Then,*

- $f_v^\perp = \left( \sum_{i=1}^k f^i \right)_v + (f_v^0)^\perp$  and  $(f_v)^\parallel = (f_v^0)^\parallel.$
- Let  $f_{v,i} := (f_v)^i$ . Then, for  $i > 0, f_{v,i} = f_v^{i+1}$  and  $f_{v,0} = f_v^1 + f_v^{0\perp}.$

*Proof.* We wish to show that for any  $i > 0$ ,  $(f^i)_v^\perp = (f^i)_v$  which is equivalent to saying that  $\langle (f^i)_v, \mathbf{1}_v \rangle = 0$  where  $\mathbf{1}_v$  is the constant function on the link  $X_v$ . Let  $I_v \in C^0$  be the function that is  $I_v(x) = 1$  iff  $x = v$ . Now, it is easy to see that  $P_{k-1}^\uparrow \cdots P_0^\uparrow I_v = \mathbf{1}_v$ . Now,

$$\begin{aligned} \langle f_v^i, \mathbf{1}_v \rangle &= \langle f_v^i, P_{k-1}^\uparrow \cdots P_0^\uparrow I_v \rangle \\ &= \langle P_0^\downarrow \cdots P_{k-1}^\downarrow f_v^i, I_v \rangle \\ &= \langle 0, I_v \rangle. \end{aligned}$$

The last equality uses the definition of an  $i$ -level cochain.  $\square$

**Theorem 2.15** (Decomposition Theorem). *Let  $f \in C^k$  and decompose  $f = \sum_i f^i$  where  $f^i \in C_i^k$ . Then,*

$$\begin{aligned} \langle P_k^\vee f, f \rangle &\leq \sum_{i=0}^k \lambda_{\phi, i, k} \|f^i\|^2 + \underbrace{\sum_{i \neq j} c_{ij} \langle f^i, f^j \rangle}_{\text{Mixed Terms (MT)}} \\ &\leq \lambda_{\phi, 0, k} \|f\|^2 + \text{MT}. \end{aligned}$$

where  $\lambda_{\phi, i, k} = 1 - \frac{1}{k+1-i} \prod_{j=i-1}^{k-1} (1 - \lambda_j)$  and  $\lambda_j = \max_{\tau \in X(j)} \lambda_2(X_\tau)$ .

**Proof sketch:** The proof is quite similar to that of the Trickle Down Theorem from the previous lecture wherein we decompose the inner product as a sum over the perpendicular and parallel components of the local cochains  $f_v$ . The perpendicular part is easily bounded (as earlier) using the inductive hypothesis on the links  $X_v$ . We make the following notation to make the induction easier.

$$\lambda_{v, i, k} := 1 - \frac{1}{k+1-i} \prod_{j=i-1}^{k-1} (1 - \lambda_{v, j}) \text{ and } \lambda_{v, j} = \max_{\tau \in X_v(j)} \lambda_2(X_\tau)$$

The parallel part is handled using the ‘‘advantage lemma’’ [Lemma 2.16](#) which can be seen as the technical core of the result.

*Proof.* We proceed by induction on  $k$ .

**Base Case:** When  $k = 0$ ,  $\langle \tilde{P}_0^\wedge f, f \rangle \leq \lambda_{-1} \|f\|^2 = \lambda_{\phi, 0, 0} \|f\|^2$ .

**Inductive Case:**

$$\begin{aligned} \langle \tilde{P}_k^\wedge f, f \rangle &= \mathbb{E}_{v \in X(0)} \left[ \langle \tilde{P}_{k-1, v}^\wedge f_v, f_v \rangle \right] && \text{Localization Lemma} \\ &= \mathbb{E}_{v \in X(0)} \left[ \langle \tilde{P}_{k-1, v}^\wedge f_v^\perp, f_v^\perp \rangle \right] + \mathbb{E}_{v \in X(0)} \left[ \langle \tilde{P}_{k-1, v}^\wedge f_v^\parallel, f_v^\parallel \rangle \right] \\ &= \underbrace{\mathbb{E}_{v \in X(0)} \left[ \langle \tilde{P}_{k-1, v}^\wedge f_v^\perp, f_v^\perp \rangle \right]}_A + \underbrace{\mathbb{E}_{v \in X(0)} \left[ \langle f_v^{0\parallel}, f_v^{0\parallel} \rangle \right]}_B. && \text{Using Claim 2.14} \end{aligned}$$

Since  $f_v^\perp$  is a cochain in  $C^{k-1}(X_v, \mathbb{R})$ , we can use the inductive hypotheses to obtain,

$$\left\langle \widetilde{\mathbf{P}}_{k-1,v}^\wedge f_v^\perp, f_v^\perp \right\rangle \leq \sum_{j=0}^{k-1} \lambda_{v,j,k} \left\| (f_v^\perp)^j \right\|^2 + \text{MT}$$

From [Claim 2.14](#), we know that for  $j > 0$ , the  $j$ -level cochain of  $f_v^\perp$  are the same the localization of the  $j + 1$ -level cochain of  $f_v$ , i.e., is  $(f_v^\perp)^j = (f^{j+1})_v$ . Plugging this in,

$$\sum_{j=0}^{k-1} \lambda_{v,j,k} \left\| (f_v^\perp)^j \right\|^2 = \lambda_{v,0,k} \left\| (f_v^0)^\perp + f_v^1 \right\|^2 + \sum_{j=1}^{k-1} \lambda_{v,j,k} \left\| f_v^{j+1} \right\|^2 \quad (2.2)$$

$$\leq \lambda_{v,0,k} \left\| f_v^0 \right\|^2 + \sum_{j=0}^{k-1} \lambda_{v,j,k} \left\| f_v^{j+1} \right\|^2. \quad (2.3)$$

Now we will use the observation that  $\lambda_{\phi,i,k} \geq \lambda_{v,i-1,j}$  for any  $v \in X(0)$  and  $j \leq k$ . Moreover, by localization lemma ([Lemma 2.12](#)), we get  $\mathbb{E}_v \left[ \left\| f_v^j \right\|^2 \right] = \left\| f^j \right\|^2$ . The term (A) then can be bounded as

$$(A) \leq \lambda_{\phi,1,k} \left\| f^0 \right\|^2 + \sum_{j=1}^{k-1} \lambda_{\phi,j,k} \left\| f^j \right\|^2.$$

□

**Lemma 2.16** (Advantage lemma).

$$\mathbb{E}_{v \in X(0)} \left[ \left\| f_v^0 \right\|^2 \right] \leq \left( 1 - \frac{k}{k+1} (1 - \lambda_{-1}) \right) \left\| f^0 \right\|^2.$$

*Proof.* Proof not provided in the lecture. See [[GK22](#), Lemma 7.8].

□

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